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A benchmarking methodology for evaluating software switch performance for NFV

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Abstract—Interest in software networking has grown significantly since the introduction of Network Function Virtualization (NFV). Software switches are used in NFV to steer traffic between different virtualized network functions and physical Network Interface Cards (NICs). It is becoming more and more important to objectively evaluate and compare the performance of the multiple alternative implementations that have recently been proposed. A comprehensive performance analysis is still missing for two main reasons: (i) the amount of time required to configure and compare all such tools is enormous; (ii) it is very difficult to define a proper methodology to compare different solutions in a fair manner. In this paper, we propose a methodology based on four simple yet representative test scenarios used to evaluate the performance of software switches. We apply this methodology to measure throughput and latency metrics for 6 state-of-the-art software switches namely, OVS-DPDK, snabb, BESS, FastClick, VPP and netmap VALE. Our work constitutes a first step to building a better understanding of design tradeoffs and identifying performance bottlenecks.

I. CONTEXT

Network Function Virtualization (NFV) is expected to have a significant impact on the networking infrastructure, considerably reducing CapEx and OpEx by replacing proprietary, expensive and inflexible hardware middleboxes with virtualized network functions (VNFs) implemented on commodity servers. Software switches are widely adopted by NFV platforms as the dataplane to steer traffic and deliver services. For example, E2 [1] and ParaBox [2] use BESS [3] as their dataplane while ClickOS [4] and HyperNF [5] adopt netmap’s VALE switch [6]. Other implementations such as OVS-DPDK [7], snabb [8], and FD.io VPP [9] also aim to deliver high performance NFV solutions. It is thus of paramount importance for network operators and service providers to understand the performance baselines and to identify the bottlenecks of software switches.

In agreement with the recent observations by Fang et al. [10], we believe that evaluating the performance of software switches is in practice a very difficult task. As noted in [10]. most works consider the evaluation of a single design at a time [11, 12, 13], or compare just of a small subset of designs in a narrow test scenario [14, 15, 16].

We have additionally observed that the metrics chosen to compare designs are not always sufficient to fully determine which one is “best”. For instance, we reproduced a very simple scenario with a software switch deployed as a L2 forwarder between two NICs, and measured overall throughput (bidirectional), latency (round-trip time in μs), and the standard deviation of the latter. We tried to understand if these metrics are correlated, or otherwise stated: is the design with the highest throughput also offering the smallest latency? Fig. 1 reports our findings. In the left plot, we report measured throughput against the latency. It may be seen that there is indeed a negative correlation, implying the switch with highest throughput is also the one providing the minimal latency. We then plotted the measured latency against the standard deviation in the measurements: this time no pattern could be detected, meaning that there is no correlation between the packet processing rate and performance stability. For instance, while OVS-DPDK can offer a high processing throughput and low latency, it does so with a very high variance in the latency measurements, thus showing that performance can be rather unstable.

We conclude that it is necessary to provide: (i) a fair comparison for all the state-of-the-art software switches, since existing works consider only one or a small subset of them; (ii) simple yet representative test scenarios that can help us to develop a better understanding of the performance baseline.

Fig. 1: Scatter plot of latency/throughput and latency/std. deviation. The throughput results shown in the left plot are obtained with bidirectional tests between two 10Gbps NICs using 64B packets.

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and identify the bottleneck. The latter is especially important given the fact that network functions are deployed in virtual environments such as virtual machines (VM) and containers.

II. METHODOLOGY

In our work, we consider 6 state-of-the-art software switch implementations, namely OVS-DPDK [7], BESS [3], snabb [8], netmap VALE [6], FastClick [17] and FD.io VPP [9]. All the selected designs are open-source projects with high performance on our testbed. In particular, OVS-DPDK is an OpenFlow (OF) switch performing match-action instructions on a per-packet basis. It can also forward packets to a remote OF controller for further processing. BESS, FastClick, Snabb and FD.io VPP are modular tools. They are capable of composing complex services and interconnecting VNFs with the network. Netmap’s VALE is a pure layer-2 switch. Other open-source software switch implementations are excluded either because of very low performance (eg., Lagopus [18]) or high degree of similarity (eg., ClickNF [19]).

We assume VNFs can run either inside virtual machines or containers and we consider throughput and latency (in terms of round-trip time) as the main performance indicators. Each switch implementation, the device under test (DUT), is evaluated with 4 different test scenarios: physical NIC to physical NIC (p2p), physical NIC to VNF (p2v), VNF to VNF (v2v) and physical NIC - VNF - physical NIC or loopback (lo), as illustrated in Fig. 2. The data flows of the 4 scenarios are explained as follows:

(p2p) physical-to-physical: In this scenario, packets arrive via a physical NIC, are fetched by the DUT and then forwarded out via the other NIC. Measurements are performed on both source and sink of the packet flow, outside the DUT. This is the most basic test scenario with no VNFs running in the virtual environment and provides the baseline results for all software switches.

(v2v) virtual-to-virtual: In the v2v scenario, packets are generated by a source VNF and injected into the DUT. The DUT then forwards the packets to a sink VNF for measurement. This scenario demonstrates a DUT’s efficiency in terms of traffic forwarding between VNFs.

(lo) loopback: In the loopback scenario, packets arrive via a physical NIC and are fetched by the DUT, which forwards them to a VNF. The VNF is a lightweight cross-connect function that forwards the received packets from one virtual interface to the DUT via another virtual interface. The DUT then proceeds to forward the packets to the other NIC. This is a typical NFV scenario with a single VNF on the service chain and provides baseline reference for scenarios of multi-VNF chains.

We believe that these four test scenarios provide a baseline performance characterisation for the software switches. The p2p scenario provides the baseline result showing how efficiently software switches can process packets between physical NICs. The p2v and v2v scenarios, representing two typical dissected hops of NFV data flows, allow us to analyse and identify potential performance bottlenecks. The loopback scenario provides results for a service chain with a single VNF and can be used as a basis to predict the performance of multi-VNF chains.

III. EVALUATION PLATFORM

Experiments are conducted on a commodity server running Linux 4.8.0-41-generic distribution. The server is equipped with 2 Intel Xeon E5-2690 v3 @ 2.60GHz CPUs (each with 24 physical cores), using 32k/256k/30720K L1-3 caches, and 2 Intel 82599ES dual-port 10-Gbps NICs. All the VNFs are hosted by QEMU/KVM virtual machines. We use MoonGen [20] as traffic generator, FlowWatcher-DPDK [21] as VNF for p2v and v2v scenarios to measure the throughput. DPDK l2fwd [22] is the VNF used to forward packets in the loopback scenario. For the packet exchange between VMs and software switches, we adopt the netmap passthrough [23] exclusively for the netmap VALE switch, and the vhost-user protocol [24] for the other considered software switches. Both enable zero-copy packet delivery through shared memory. For all the experiments, we use identical settings for the virtual machines and software switches to ensure a fair comparison. We pin threads of the software switches to isolated CPU cores and disable Turbo Boost to reduce performance variance.

1 According to our benchmarking result, Lagopus cannot achieve more than 2 Mpps even for the most simple scenario.
2 ClickNF, just like FastClick, is an extension of the Click Modular Router. Our benchmarking result shows ClickNF achieves similar throughput and latency with respect to FastClick.
In this section, we provide some preliminary results for the throughput test. Fig. 3 illustrates the unidirectional throughput for the considered software switches in the scenarios described in Fig. 2. As we observe, no single software switch prevails in the considered scenarios. For example, BESS outperforms the others in the p2v case thanks to its efficient implementation of packet delivery to virtual machines. Snabb prevails in the v2v scenario because its vhost-user implementation is more efficient for packet delivery between virtual machines.

IV. EXPERIMENTAL RESULTS

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