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A comprehensive approach to the automatic refinement and verification of access control policies

Manuel Cheminod\textsuperscript{b}, Luca Durante\textsuperscript{b}, Lucia Seno\textsuperscript{b}, Fulvio Valenza\textsuperscript{a,b,*}, Adriano Valenzano\textsuperscript{b}

\textsuperscript{a}Dipartimento di Automatica e Informatica, Politecnico di Torino, Torino, Italy
\textsuperscript{b}CNR-IEIIT, c.so Duca degli Abruzzi 24, Torino I-10129, Italy

Abstract

Access control is one of the building blocks of network security and is often managed by network administrators through the definition of sets of high-level policies meant to regulate network behavior (policy-based management). In this scenario, policy refinement and verification are important processes that have to be dealt with carefully, possibly relaying on computer-aided automated software tools.

This paper presents a comprehensive approach for access control policy refinement, verification and, in case errors are detected in the policy implementation, their fixing. The proposed methodology is based on a twofold model able to describe both policies and system configurations and allows, by suitably processing the model, to either propose a system configuration that correctly enforces the policies, or determine whether a specific implementation matches the policy specification also providing hints on how possible anomalies can be fixed. Results on the average complexity of the solution confirm its feasibility in terms of computation time, even for complex networked systems consisting of several hundred nodes.

Keywords: Access Control, Policy-Based Network Management, Policy Refinement, Policy Verification.

1. Introduction

Access control is a major building block of network security\textsuperscript{[1]}, regulating access of legitimate users to resources of a system. Access control by itself cannot prevent the occurrence of cyber-attacks, however effective protection schemes cannot abstract from both the definition of how access to system resources shall take place as well as the correct enforcement of the desired behavior in the system.

As target networked systems can be very complex, administrators often rely on policy-based paradigms for access control management (i.e., policy-based management, PBM\textsuperscript{[2]}). Policies are technology-independent rules which define the desired network behavior from a high-level perspective. As such, they allow to separate the two problems of specification (i.e., definition of the desired network behavior) and implementation (i.e., actual enforcement of the desired behavior in the system) making network management easier and more flexible.

Access control policies regulate the access of users to the resources of a networked system by defining "who is allowed to do what on what". The application of PBM to access control has been receiving increasing attention by the scientific community since many years, and research on this topic has been focusing on three main domains, namely policy analysis (PA), policy verification (PV) and policy refinement (PR).
Policy analysis deals with the fulfillment of specific properties by a set of policies \cite{3}. The goal of PA is the detection of specification inconsistencies, arising when two or more policies are conflicting. PA only concerns the specification domain (i.e., the policies) and does not consider the implementation domain, i.e., how policies are actually enforced in the system.

Policy refinement bridges the gap between specification and implementation. Indeed, a PR process translates high-level policies into low-level system configurations \cite{4}. PR is a critical task that, if not carefully performed, may lead to either incorrect or sub-optimal implementations, thus affecting important network performance indexes such as throughput and even jeopardizing the overall system security. When PR is performed by hand the risk of errors increases. The Verizon Data Breach Investigation Report points out that 14\% of breaches in 2017 were caused by human errors in network configuration \cite{5}. For this reason, computer-aided automatic or semi-automatic tools are definitely needed to assist administrators in the translation of access control policies.

Finally, PV deals with checking whether a set of policies is correctly enforced in a system or, in other words, whether the system implementation actually matches the policy semantics and no anomalies are present. As an example, policies may state that “administrators can login in all nodes”, but a mismatched firewall configuration may result in cutting off some server. PV can also be used to validate a hand-made refinement. Note that, sometimes, performing a full refinement of access control policies is not convenient. In particular, when policies change in large networks, administrators verify whether the new policies are already correctly enforced with the current system configuration and, only if the result is negative, they introduce the required modifications. In this case the verification process plays a crucial role.

In this paper we present a comprehensive approach to automatic access control (AC) policy refinement and verification. Our proposal builds on both the twofold model introduced in \cite{6}, which enables the description of high-level AC policies and low-level implementation details of a target networked system, and some promising results described in \cite{7}.

While the solution in \cite{6} was only meant for a posteriori policy verification, the technique shown in this paper extends \cite{6} by focusing on credential assignment as a means to enforce AC policies. This is leveraged to solve the following problems, commonly tackled by security administrators:

1. Is there a user credential assignment which allows the implemented system to satisfy the policies? (policy refinement);
2. Can a partial user credential assignment be completed so that the system behavior fulfills the policy specification? (constrained policy refinement);
3. Does a specific user credential assignment make the system behavior match the policy specification? (policy verification);
4. How should the credential assignment be modified to fix anomalies (i.e., discrepancies between specification and implementation)? (anomaly fixing).

In practice, questions above refer to three different situations, that is:

1. no credential has been granted yet;
2. some credentials have already been assigned (constrained credentials);
3. all credentials have been assigned.

Rough connections to the state of a system development are then: under design, partially implemented and fully implemented.

The paper is structured as follows: Sect. \ref{sec:overview} summarizes the characteristics of the twofold model employed
for policy and system description. Sect. 3 and 4 respectively, formally present the proposed approach (and how to solve the considered problems) and provide conceptual workflows describing how administrators can be assisted in managing access control for their network in different circumstances. Sect. 5 describes some optimizations introduced in the actual implementation of the proposed solution and Sect. 6 deals with the automated tool and its performance. Finally, Sect. 7 considers some related works and Sect. 8 draws some conclusions and provides hints on future work.

2. Background

The twofold model presented in §2 enables the description of both a set of high-level access control policies (through a Specification Model $S$) and the fine-grained details of a target networked system (through an Implementation Model $I$). Model $I$, in particular, stores information about the system structure, components and configurations, i.e., nodes and their interconnections, communication protocols, physical places where nodes are placed (e.g., rooms and cabinets), doors between adjacent spaces, software applications running on nodes, available and exploited access control mechanisms, credentials owned by users, etc.

By suitably processing models $S$ and $I$, three sets of pairs, $S^+$, $S^-$ and $I$ can be computed such that

$$S^+, S^-, I \subseteq Users \times Actions \tag{1}$$

where $Users$ and $Actions$ are sets identifying, respectively, all users $u$ possibly interacting with the system and all actions $\alpha$ that can be performed in the system. In the following, we use $u$ and $\alpha$ to mean a generic element of set $Users$ and $Actions$, respectively.

Sets $S^+$ and $S^-$ derive from $S$ and define, respectively, allowed and denied users actions. Basically, whenever a pair $(u, \alpha) \in S^+$ it means that, according to the policies, user $u$ is allowed to perform action $\alpha$ on the system. Conversely, if $(u, \alpha) \in S^-$, according to the policies, user $u$ must not be able to perform action $\alpha$ on the system.

In the following, $S^+_u$ and $S^-_u$ are the sets of allowed and denied actions specified for user $u \in Users$, i.e.

$$S^+_u := \{(u, \alpha) \in S^+ | u = \tilde{u}\} \tag{2}$$

Clearly, $S^+ = \bigcup_{\tilde{u} \in Users} S^+_\tilde{u}$. Moreover, we assume $S^+ \cap S^- = \emptyset$ and, consequently, $S^-_u \cap S^+_u = \emptyset \ \forall u \in Users$, (i.e., policies are supposed to be coherent and, as such, conflict free).

The Implementation set $I$ is derived from $I$, and pairs belonging to $I$ describe all actions that users are actually allowed to perform on the system, as determined by its current implementation.

Set $I$ is only a portion of the information obtained by processing model $I$. Indeed, by suitably combining $I$ with a set of purposely defined inference rules, a collection of automata $A_u$ can be computed (one automaton for each user $u \in Users$), such that

$$A_u = (Q_u, \Sigma_u, \delta_u, q_u^0) \tag{3}$$

Equation (3) adopts the traditional notation for Deterministic Finite State Automata (DFSA) §5 where $Q_u$ is the set of states $q_u$, $q_u^0 \in Q_u$ is the initial state, $\Sigma_u \subseteq Actions$ is the set of events, i.e. user's actions, and $\delta_u$ is the transition function of the automaton. Overall each automaton $A_u$ stores the sequences of actions $\alpha$ allowed to user $u$, given a specific system implementation described by $I$, and how her/his state varies accordingly.

Inference rules describe the way user/system interactions take place (e.g., how a user can move, stepping a door, between adjacent rooms or how s/he can access a file if s/he is logged on its host and has the required privi-
leges and/or credentials). Each automaton $A_u$ is computed by considering the interaction of user $u$ with the system and applying all inference rules to register actions s/he is able to perform and how this affects her/his state (i.e., the room s/he is in, access to local resources s/he gains).

Automata $A_u$ are used to derive sets $I_u$. Each $I_u$ consists of all pairs $(u, a)$ obtained by combining user $u$ with all $a$ appearing in at least one transition of $A_u$. The implementation set $I$ can then be obtained as

$$I = \bigcup_{u \in Users} I_u$$

The problem of policy verification can formally be assessed by comparing sets $S^+$, $S^-$ and $I$. In particular, by defining the anomaly sets containing all actions respectively allowed and denied to users in violation of the policies as

$$\bar{S}^+ := S^+ \setminus I$$
$$\bar{S}^- := S^- \cap I$$

we obtain that the implementation satisfies the policy specification if and only if

$$\bar{S}^+ = \emptyset \quad \land \quad \bar{S}^- = \emptyset$$

Note that for the implementation to match the specification, $I = S^+$ is not a necessary condition. Indeed, the high-level specification of policies may put into evidence only a subset of actions that are possible in the systems. In other words $I$ may contain pairs not explicitly considered in the policies defined by the network administrators. As an example, $I$ may contain a pair $(u, a)$ where $a = (\text{login, PC})$. This kind of actions is useful to describe how access to some of the system resources actually takes place (e.g., a user, for accessing a specific software application running on PC, may first need to login on the PC itself). Often, these actions are not taken into account by policies, which in turn, do not have any knowledge of where applications are actually installed. Policies may explicitly state that $u$ shall be able to access a specific software application without mentioning logging in on the PC hosting the application itself. If not otherwise specified in the denied action set, this means that for a correct implementation it is not important whether user $u$ is able to login on PC or not.

Since automata are built on a per-user basis, an easy way to check the correctness of the policy implementation is by computing (4) and checking (5) $\forall u \in Users$ individually, i.e.:

$$\bar{S}^+_u := S^+_u \setminus I_u$$
$$\bar{S}^-_u := S^-_u \cap I_u$$

$$\bar{S}^+_u = \emptyset \quad \land \quad \bar{S}^-_u = \emptyset$$

3. Formal description of the approach

Approach in [6] only focused on policy verification and, in particular, on computing sets $\bar{S}^+$ and $\bar{S}^-$ and checking (5), to possibly highlight the need for changes in the current implementation. However, such a technique was unable to provide hints on how these system modifications had to be carried out, leaving this critical task completely to the network administrators.

This work extends the scope of [6] by including both policy refinement and verification in the proposed solution and by automatically suggesting fixings for the detected anomalies. Modifications to the implementation (reflecting on $I$), may, in principle, involve different aspects of the target system (e.g., network topology, node configurations, etc.). Of course, some changes are very complex and may affect the system functionality. In this paper we take into account only changes pertaining to user credential assignments, as they are easy to manage and often help with fixing a large number of detected anomalies.

The main idea is determining relations among actions $a$ and credentials needed to perform them, i.e., what sequences of actions and, consequently, what sets of creden-
tials allow users to gain privileges enabling the execution of a specific action. This information is already present in $I$, as it is necessary to check whether an inference rule can fire, leading to a state transition in $G$. In the following $C$ is the set of all credentials in $I$, and set $C_u \subseteq C$ if $u \in \text{Users}$ stores the credentials owned by user $u$. When a transition can fire in $A_u$, the user’s credential $c \in C_u$ allowing the corresponding action is known, even though it was not included in the transition label in $G$ because of its restricted goal.

By contrast, credentials are also used here to fix anomalies detected through sets $G$, and this imposes managing $c$ allowing $a$ explicitly in the corresponding automaton transition labels. More formally, from now on, each transition of any automaton $A_u$ is associated with an event $e \in \Sigma_u \subseteq \text{Actions} \times (C \cup \{\varepsilon\})$. Thus, in general, event $e$ has the form $(a, c)$, i.e., it is a pair combining an action and a credential requested to perform that action (if any). Symbol $\varepsilon$ means “no credential”, i.e., whenever a transition is labeled with $e = (a, \varepsilon)$, no specific credential is necessary to carry out $a$. All algorithms presented in $G$ to build automata $A_u$ still work properly here, considering events as described above.

To minimize the computational effort, we compute a single automaton $A_u$ referring to a hypothetical super-user $u = \text{sup}$. Superuser $\text{sup}$ is defined as the user who is assigned all credentials existing in the system, i.e., $C_{\text{sup}} = C$. Automaton $A_{\text{sup}} = (Q_{\text{sup}}, \Sigma_{\text{sup}}, \delta_{\text{sup}}, q^0_{\text{sup}})$ is built exactly in the same way as automata $A_u$. While each automaton $A_u$ describes the interactions between the system and a specific user $u$, automaton $A_{\text{sup}}$ represents the most powerful interaction allowed by the system characteristics. Indeed, as it models the interaction achievable by owning all possible credentials, $A_{\text{sup}}$ generates a maximum language consisting of all sequences of actions any user can perform in the system. As such, $A_{\text{sup}}$ describes the system global behavior and is useful to examine dependencies among actions avoiding the computation of as many $A_u$s as $u \in U$. Information about any user $u \in U$ can easily be derived directly from $A_{\text{sup}}$ whenever needed.

The following sections deal with the computation, from automaton $A_{\text{sup}}$, of expressions specifying whether a user is enabled to carry out an action depending on the credentials s/he has. Moreover, we also discuss how these expressions can help in tackling the refinement, constrained refinement and verification problems. In more detail, starting with a generic DFSA, section 3.1 shows how to compute a collection of sets of enabling events $\forall e \in \Sigma$, such that $e$ can be impeded by disabling at least one event in each set. Section 3.2 instead, makes use of sets of enabling events for automaton $A_{\text{sup}}$ to determine $\forall e \in \Sigma_{\text{sup}}$ a boolean function (enabling function) expressing what credentials are needed to perform the action guarded by $e$. Finally, sections 3.3, 3.4 and 3.5 apply enabling functions to policy refinement, constrained refinement and verification respectively.

3.1. Enabling events

Let us consider the simple automaton $G = (Q, \Sigma, \delta, q_0)$ shown in Fig. 1 where events $e \in \Sigma$ are $(a, c)$ pairs mentioned above. $G$ is a super-user automaton computed for an example system where:

- two adjacent rooms (A and B) are connected through a door which can be opened with key $k_{AB}$;
- two system nodes (host $H_1$ and server $S_1$), connected through the network infrastructure, are located in $B$;
- two different passwords (user and admin) can be used to sign in on $H_1$;
- whenever a user is logged in as an administrator on $H_1$, s/he can perform the backup action (e.g., invoke a backup service), while simple users are not allowed to do so;
- any user logged in $H_1$ can access, remotely through the network, a database hosted by $S_1$. 

For sake of simplicity short labels have been used for edges in $G$, i.e., $\Sigma = \{a, b, c, d, x, y\}$, while corresponding (complete) labels obtained through the computation of the automaton from $I$ are listed in the table below the automaton, e.g., event $a$ stands for a pair $(a, c)$ where action $a$ is $\text{enter}$, $A$ whereas credential $c$ is $k_{AB}$.

**Definition 1.** Given an event set $\Sigma$ and a string $s \in \Sigma^*$, we define operator $E(\cdot) : \Sigma^* \rightarrow 2^\Sigma$ as:

$$E(s) = \{e \in \Sigma \mid \exists t, v \in \Sigma^* : s = tev\}$$

Given a string $s$ of events of $\Sigma$, operator $E(\cdot)$ returns the set of all events in $s$. For instance, if we consider strings $s_1 = \text{acca}by$, $s_2 = abu$ and $s_3 = e$ belonging to $\mathcal{L}(G)$ in Fig. 1 and apply $E(\cdot)$ to each one of them, we obtain $E(s_1) = \{a, c, y, b\}$, $E(s_2) = \{a, b\}$ and $E(s_3) = \{e\}$, respectively.

**Definition 2.** Given an automaton $G = (Q, \Sigma, \delta, q_0)$ and an event $e \in \Sigma$, a set of events $E \subseteq \Sigma \setminus \{e\}$ enables event $e$ in $G$ ($E \triangleright e$) as:

$$E \triangleright e = \begin{cases} \exists s \in \Sigma^* : \\ & e \in \mathcal{L}(G) \\ & E(s) = E \\ & E(s') \subseteq E \end{cases}$$

Basically, a set of enabling events for an event $e$ is the smallest set of events whose combined occurrence enables event $e$.

For instance, considering $G$ in Fig. 1 we have $\{a\} \triangleright b$, $\{a\} \triangleright c$, $\{a\} \triangleright d$, and $\{a, d\} \triangleright y$, while $\{a, d, x\} \not\triangleright y$ being $\{a, d\} \subseteq (a, d, x)$.

**Definition 3.** Given an automaton $G = (Q, \Sigma, \delta, q_0)$ and an event $e \in \Sigma$, we define operator $En(\cdot) : \Sigma \rightarrow 2^\Sigma$ as:

$$En(e) = \{E \subseteq \Sigma \mid E \triangleright e \text{ in } G\}$$

Given a generic automaton modeling a system, any event $e$ in its set $\Sigma$ is characterized by a (possibly empty) set of sets of enabling events $En(e)$. From the definitions above it follows that for $e$ to occur, all events in at least one of sets $E \in En(e)$ must happen and, consequently, to prevent $e$ from happening at least one element in each set $E \in En(e)$ must be disabled. Moreover, from (8), any pair of sets $E_1, E_2 \in En(e)$ may share some elements but $E_1 \cap E_2$ is never true. Note that the set of enabling events for $e \in \Sigma$ may be $En(e) = \{\{e\}\}$ meaning that $e$ can only occur directly in the initial state of $G$ (i.e., it is not enabled by other events). Moreover, $En(e) = \{\{e\}\}$ is different from $En(e) = \emptyset$, as the latter means that $e \in \Sigma$ can never happen in $G$ (i.e., no state having an outgoing transition labeled as $e$ is reachable from $q_0$).

The computation of $En(e), \forall e \in \Sigma$ for $G$ in Fig. 1 leads
3.2. Enabling functions

The following definitions apply to automata having the characteristics of $A_{sup}$. In particular, given an event $e = (a, c) \in \Sigma$, we define $C(\cdot) : \Sigma \rightarrow (C \cup \{e\})$ and $A(\cdot) : \Sigma \rightarrow \{\text{Actions}\}$ such that:

\[
C(e) = c \quad \text{and} \quad A(e) = a
\]

Each event $e = (a, c)$, can be disabled for user $u$ either directly (by denying credential $c$ to $u$), or indirectly (by disabling at least one event in any set belonging to $E \in En(e)$ for $u$). Note that events $e$ where $C(e) = e$ can be disabled only indirectly.

In this condition, an event $e$ can happen only if a user owns credential $c$ for $e$ together with the credentials for all events in at least one of the sets in $En(e)$. This can be described by means of a boolean expression where a variable is associated to each credential, whenever the value of the variable is 0 it means that the credential is not provided to the user, while, on the contrary, when the variable assumes value 1 the user is assigned the corresponding credential. A sum of variable products is then specified for each $E \in En(e)$.

As an example, let us consider event $y$ in Fig. 1: $En(y) = \{(a, c), (a, d)\}$. The corresponding action $(\text{access, DB})$ can be performed only by users who own the credential for $y$ together with those allowing $(a$ and $c)$ or $(a$ and $d)$, i.e.

\[
pw_\text{db} \cdot k_{AB} \cdot pw_\text{uh} + pw_\text{db} \cdot k_{AB} \cdot pw_\text{ah} = \text{true} \quad (10)
\]

Any assignment of variables making expression (10) true corresponds to a set of credentials enabling $u$ to perform action $(\text{access, DB})$ in the example.

More in general, for each $e = (a, c) \in \Sigma_{sup}$, given its set of sets of enabling events $En(e)$, we build a boolean function $F(a)$ (enabling function) as sketched above.

**Definition 4.** Given the set $C$ of all credentials in $I$, we define a set of boolean variables

\[
V = \{v_c \mid c \in C\} \quad (11)
\]

where $v_c$ is the variable associated to credential $c$.

For the example automaton in Fig. 1 we have that set $V = \{pw_\text{db}, kw, pw_\text{uh}, pw_\text{ah}\}$, where each variable has been named according to the associated credential.

**Definition 5.** Given an automaton $A_{sup} = (Q_{sup}, \Sigma_{sup}, \delta_{sup}, q_{sup}^0)$ and an event $e \in \Sigma_{sup}$, $f(e)$ is the boolean expression

\[
f(e) := v_c \cdot \left( \sum_{E \in En(e)} \prod_{e \in E} v_{C(e)} \right) \quad (12)
\]

For instance, $f(y)$ for system $G$ in Fig. 1 is equation (10).

Event $e = (a, c)$ can occur only if the user has credential $c$, i.e. $f(e)$ cannot be true if $v_c$ is false. In addition, the user must own all credentials requested for all events in at least one $E \in En(e)$, i.e. the corresponding variable product must be true.

In the example system $G$, action $(\text{login}, H_1)$ appears in both events $c$ and $d$, and credentials $pw_\text{uh}$ and $pw_\text{ah}$ are respectively needed to enable them. This situation, which happens frequently in real systems, has to be handled explicitly as policies deal with actions irrespectively of the credentials needed to carry them out.

Definition 5 can then be extended to define how action $a$ can be performed whatever the credential requested for it:

**Definition 6.** Given an automaton $A_{sup}$ and an action $a$ such that $(a, c) \in \Sigma_{sup}$

\[
F(a) := \sum_{c \in C \mid (a, c) \in \Sigma_{sup}} f(a, c) \quad (13)
\]
Given an automaton

Definition 7. A

in

e
check

correct

correction.

when a complete, already available assignment has to be

administrators, while verification and (possibly) fixing is

partial credential assignment already provided by network

strained refinement is similar and consists of completing a

renement
dential assignment allowing

uxing

ties. Renement means nding a user's cre-

Basic definitions above are useful to address the policy

refinement, constrained refinement, and verication and fixing problems. Refinement means nding a user’s creden-
tial assignment allowing u to carry out all actions in

S

u,

and preventing u from performing any α ∈ S

u−.

Constrained refinement is similar and consists of completing a

partial credential assignment already provided by network

administrators, while verication and (possibly) fixing is

when a complete, already available assignment has to be

checked for correctness.

Some actions α belonging to Actions might not appear

in some e ∈ Σ

sup.

When this happens, i.e. α is not present

in A

sup.

F(α) = 0 by definition:

Definition 7. Given an automaton A

sup.

and an action

α | F(α, c) ∈ Σ

sup.

F(α) := 0

(14)

3.3. Policy renement

Our solution for policy renement is meant to determine a user credential assignment such that the system

implementation matches the policy specication, i.e., (5) holds true. The process works on a per-user basis and tries

to select values for variables of V, such that each user is

enabled to perform all allowed actions (i.e., all (u, α) ∈ S

u+)

while is prevented from performing all denied actions (i.e.,

all (u, α) ∈ S

u−).

From (13) and (14) u is able to perform all allowed

actions if and only if:

F(α) = 1 ∀ α | (u, α) ∈ S

u+

Similarly, the system prevents u from performing any of

the denied actions if and only if:

F(α) = 1 ∀ α | (u, α) ∈ S

u−

By combining conditions above, a suitable credential assignment for u is such that

(15)

Consequently we can assert that, in our case, solving the renement process means finding, for each user

u, an assignment of values for variables in V such that (15) holds.

Note that such an assignment is partial with respect to V, since values are only found for variables explicitly

appearing in (15). Variables not appearing in (15) can be safely set to 0 without affecting the policy implementa-
tion for user u. Thus a complete assignment σ

u : V → {0, 1},

is obtained by simply assigning default values to variables not present in (15).

When no solution is found, there is no way for the current system implementation I to satisfy the policies by

acting on the credential assignment only, and the problem is marked as unsat, i.e. no σ

u(·) is de ned. In this case, administrators have either to change other elements of the data model (i.e., system topology, firewall and service con-

igurations and so on) or relax some policy constraints if

this is compatible with the system security requirements.

3.4. Policy constrained renement

Sometimes administrators have already identied a partial credential assignment for user u, i.e., a function σ

u : V

u → {0, 1} such that any v

∈ V

u ⊂ V is bound to an un-

changeable value σ

u(v

). The renement process should
then complete the assignment so as to make \( (5) \) true.

Solving the constrained refinement problem means finding for each user \( u \) an assignment for \( \{ v_c \} = V \setminus V_u \), such that \( (15) \) holds when each \( v_c \in V_u \) is replaced by \( \sigma_u(v_c) \).

Even in this case, the solution satisfying \( (15) \), if any, can be used to build \( \sigma_u : V \to \{ 0,1 \} \), by joining it to \( \sigma_u^c(\cdot) \), and setting all variables which are not bound by \( \sigma_u^c(\cdot) \) and do not appear in \( (15) \) to the default value (0).

3.5. Policy verification

A complete credential assignment for user \( u \) is a function \( \sigma_u^V : V \to \{ 0,1 \} \) where each \( v_c \in V \) is bound to either 0 or 1. Thus, solving the verification problem means checking, for each user \( u \), whether a credential assignment \( \sigma_u^V(\cdot) \) enables a correct implementation of policies in the system.

In particular, sets \( S_u^+ \) and \( S_u^- \) in \( (3) \) can be computed by means of the enabling functions as:

\[
\begin{align*}
\tilde{S}_u^+ &= \{(u,a) \mid F(a)|_{\sigma_u^V} = 0 \land (u,a) \in S_u^+ \} \\
\tilde{S}_u^- &= \{(u,a) \mid F(a)|_{\sigma_u^V} = 1 \land (u,a) \in S_u^- \}
\end{align*}
\]

where notation \( F(a)|_{\sigma_u^V} \) stands for “\( F(a) \) where each variable \( v_c \) is replaced by \( \sigma_u^V(v_c) \)”.

Three additional formal definitions are used in the following sections:

**Definition 8.** Given an action \( a \), \( V_a \subseteq V \) is the set of all boolean variables \( v_c \) appearing in the enabling function \( F(a) \)

\[
V_a = \{ v_c \mid v_c \text{ is a variable in } F(a) \}
\]

**Definition 9.** Given an automaton \( A_{sup} \) and a user \( u \), \( \tilde{K}_u \) is the set of all actions:

\[
\tilde{K}_u := \{ \mathcal{A}(e) \mid e \in \Sigma_{sup}, (u, \mathcal{A}(e)) \notin \tilde{S}_u^+ \cup \tilde{S}_u^- \}
\]

In practice, \( \tilde{K}_u \) is the set of actions which, allowed or denied under the current system configuration, are not identified as anomalies.

**Definition 10.** Given \( K_u \), \( \bar{V}_u \) is the set of all variables:

\[
\bar{V}_u := \bigcup_{a \in K_u} V_a
\]

Basically, \( \bar{V}_u \) is the set of variables \( v_c \) associated to credentials enabling actions in \( K_u \).

4. Workflow

This section shows how the approach formally presented in Sect. 3 is put at work, while most significant details about the implementation of our technique are provided in Sects. 5 and 8.

From a practical point of view, we can distinguish two different operating phases, which are respectively referred to as offline and online. The offline phase (see Fig. 2) aims at modeling and analyzing the whole system and, consequently, is computation intensive and can be rather time-consuming. Fortunately, it is seldom carried out in practice, since changes in the system characteristics such as i.e., host locations, network topology, installed services etc. do not occur frequently with respect to modifications of access control policies. The online phase (see Fig. 3), instead, should be performed every time administrators configure or change some access control policies for any given user/group of users, or whenever the correctness of configurations has to be rechecked.

More specifically, the offline process (Fig. 2) is responsible for computing the set of actions whose execution enables action \( a \). Starting from \( I \) and a set of predefined inference rules automaton \( A_{sup} \) is generated and stored into
a database (automaton generator block in Fig. 2). The second step of the offline phase (functions computation) consists of visiting $A_{sup}$ to compute the enabling functions $F(\alpha)$ in (13) for all possible actions in the system.

The online phase works on a per-user basis and includes two distinct sub-processes, i.e. refinement and verification and fixing in Fig. 3.

4.1. Refinement process

Refinement aims at identifying a credential assignment $\sigma_u(\cdot)$ for any user $u$ (or group of users), as requested to correctly enforce the policies. Note that this problem may have one, many or no solution at all.

By assuming that all solutions are equivalent, a Satisfiability Modulo Theories (SMT) solver [9] is used to find either a correct assignment or a proof that no satisfactory assignment exists. In details, the refinement process (RP) tries to solve equation (15) by taking both the enabling (boolean) functions $F(\alpha)$ and the user policies ($S^+_u$ and $S^-_u$) as inputs. RP returns either a set of credentials $\sigma_u(\cdot)$ that correctly enforce the policies, or a counter-example for non-satisfiability. The counter-example is then reported to the administrators along with the set of conflicting policies. It is up to the administrators analyzing results and possibly introduce changes in either the policies or the data model (that is the system).

When dealing with constrained refinement, RP takes the administrator-defined credential assignment $\sigma'_u(\cdot)$ as additional input (dotted block in Fig. 3 and binds the relevant credentials to the predefined values. Of course, the larger the number of bounded variables the smaller the solution space and, consequently, the higher the likelihood of unsatisfiability.

Typically, the problem of constrained refinement arises when administrators directly assign some credentials for practical or business-related reasons, e.g., employees are usually directly assigned the password for accessing their workstations or administrators are assigned credentials for logging in as superusers on some hosts. In all these cases credential assignment is already partially defined.

4.2. Verification and Fixing

Given a credentials assignment $\sigma'_V(\cdot)$ for each user $u$, possible anomalies are detected by computing (16) so as to produce sets $\bar{S}^+_u$ and $\bar{S}^-_u$.

Moreover, when anomalies are detected, a possible fixing (if any) is automatically suggested to administrators. The fixing computation is based on a constrained refinement, where the credential assignment for actions not belonging to $\bar{S}^+_u \cup \bar{S}^-_u$ is left unchanged. Credentials only concerning actions in the anomaly sets (i.e., not belonging to set $\bar{V}_u$ in (19)) are left unbound, and a refinement is then searched. Roughly speaking, the fixing process (FP) tries to leave as many credentials as possible unassigned, and puts back into the game only those variables that are directly involved in anomalies, i.e. credentials enabling actions belonging to the anomaly sets. In case no solution is found, administrators can take different actions such as modifying the initial credential assignment or exploiting other characteristics of $I$. Options available to administrators when the automatic fixing fails are not in the scope
of this paper.

5. Offline phase optimization

Some properties of $A_{sup}$ can be leveraged to tackle the problem of computing $F(a)$ and to solve it efficiently.

The computation of an enabling function $F(a)$ relies on the identification of sets of enabling events $En(e)$. Finding $En(e)$, in a generic automaton $G$, means finding all those paths $p$ such that $pe \in \mathcal{L}(G)$ and $E(p) \in En(e)$.

This leads to an enumeration problem whose complexity depends on the size and structure of $G$. Unfortunately, given the number of events $|\Sigma|$, the number of paths to the target event $e$ can be as large as $|\Sigma|!$. Actually, the worst case condition occurs when $e$ is enabled by a combination of all the other events disjointedly. The computation complexity in the worst case is $O(|\Sigma|!)$, so an optimization strategy is needed to make the problem manageable. The main idea we followed is splitting the problem into two simpler parts (divide et impera approach), solve them individually and then recombine the partial results to efficiently obtain $En(e)$ for each $e \in \Sigma$.

5.1. Divide et impera

Given its characteristics, $A_{sup}$ can be split into two simpler automata $A'$ and $A^L$ ($A' \Leftrightarrow A_{sup}$ and $A^L \Leftrightarrow A_{sup}$), where $A' = (Q', \Sigma', \delta', q_0')$ and $A^L = (Q^L, \Sigma^L, \delta^L, q_0^L)$. Their combination by means of the standard parallel operation $\boxplus$, returns $A_{sup}$ ($A'||A^L = A_{sup}$). This enables the computation of two enabling sets of events for $A'$ and $A^L$ (respectively $En'(e)$ and $En^L(e)$ in the following) which can be combined to build $En(e)$.

Automata $A'$ and $A^L$ are built by considering complementary aspects of the system. Indeed, $A'$ takes into account both the physical interactions of a user with the nodes in the system and her/his movements between rooms.

Instead $A^L$ only deals with aspects concerning the acquisition and exploitation of logical accesses to the devices/services. In building $A^L$ we assume that the user can reach, moving from the initial state, all locations that have been explored while building $A'$ (see Axiom 1 in the following).

In general, actions $a \in Actions$ belong to three disjoint groups: $\theta_1$, $\theta_{II}$ and $\theta_{III}$. Actions in $\theta_1$ depend on the user’s physical position and can change it. Actions in $\theta_{II}$ depend on the user’s position and, possibly, let her/him gain access to a device. Finally, actions in $\theta_{III}$ depend on the set of logical accesses already acquired by the user and, possibly, let her/him enlarge it with a new element.

The construction of $A'$ combines only $a$ in $\theta_1$ and $\theta_{II}$, while $A^L$ takes into account $\theta_{II}$ and $\theta_{III}$. Events that are in both $\Sigma'$ and $\Sigma^L$ belong to $\theta_{II}$, and $A'$ describes how sequences of $\theta_1$ actions can enable a subset of $\theta_{II}$. Similarly, $A^L$ shows how the subset of $\theta_{II}$ can trigger sequences in $\theta_{III}$.

For example, $G$ in Fig. 1 can be obtained by the parallel composition of the two automata $A'$ and $A^L$ shown in Fig. 4 and Fig. 5 respectively. In this case $Actions = \{a, b, c, d, x, y\}$, $\theta_1 = \{a, b\}$, $\theta_{II} = \{c, d\}$ and $\theta_{III} = \{x, y\}$.

Let $\Gamma(q)$ be the set of events enabled in state $q$: $\Gamma(q) = \{e \in \Sigma \mid \delta(q, e) \text{ is defined}\}$. For instance, $\Gamma(q_0') = \{c, d\}$ and $\Gamma(q_1') = \{c, y, d\}$ in Fig. 5.

The complete and formal definition of the procedures used to build $A$, $A'$ and $A^L$ has been provided in [6]. Some useful properties, referred in the following as axioms, stem from the characteristics of the data model and the inference rules. Axioms, allow to derive three new properties needed to build $En(e)$ from $En'(\cdot)$ and $En^L(\cdot)$ in a very

---

2Since a string $s \in \Sigma^*$ is a sequence of events, every path $p$ in $G$ corresponds to a string in $\mathcal{L}(G)$ and we may refer to $p$ as either a path or a string indifferently.
efficient way.

**Axiom 1.** All actions $a \in \theta_{II}$ that are allowed in $A'$ belong to $A^L$ as well and, in particular, are enabled in the $A^L$ initial state $q_0^L$.

$$\forall e \in \Sigma' \cap \Sigma^L, e \in \Gamma(q_0^L), En^L(e) = \{\{e\}\}$$

**Axiom 2.** Given a state $q \in Q^L$, each event $e \in \Gamma(q)$ is also enabled in any state $q'$ reachable from $q$ ($e \in \Gamma(q')$).

**Axiom 3.** Any action in $\theta_{II}$ is enabled by a single logical access (and not by a simultaneous combination of multiple logical accesses acquired by the user).

From Axiom 2 and Axiom 3 we have:

**Property 1.** Given any two events $e_1, e_2$ in some set $E$ in the enabling set $En(e)$ of a target event $e$, either $e_1$ is in the enabling set of $e_2$ or $e_2$ is in the enabling set of $e_1$.

$$E \in En(e), \{e_1, e_2\} \subseteq E \Rightarrow \exists E' \in En(e_2), e_1 \in E' \lor \exists E' \in En(e_1), e_2 \in E'$$

From Axiom 1 and Property 1 we have:

**Property 2.** For any event $e$ in $\Sigma^L$ and any $E \in En^L(e)$ two distinct events $\{e_1, e_2\} \subseteq E$ both belonging to $\Sigma' \cap \Sigma^L$ can never exist.

Moreover, in building $A^L$ we exploit only actions belonging to $\theta_{II}$ and $\theta_{II}$. Remembering that $\theta_{II}$ events only depend on logical accesses conditions already gained by the user, and the initial set of such conditions for any user is empty, any logical access in $A^L$ can be acquired only by means of a $\theta_{II}$ action, that is events in $\Sigma' \cap \Sigma^L$. As such, any sequence of events leading to $e \in \Sigma^L \setminus \Sigma'$ starts with some specific action $\tilde{e} \in \Sigma' \cap \Sigma^L$. Since all $e \in \Sigma' \cap \Sigma^L$ are included in $\Gamma(q_0^L)$, we can conclude that:

**Property 3.** For any event $e$ in $A'$, there is exactly one event $\tilde{e}$ in each $E \in En^L(e)$, such that it also belongs to $\Sigma'$:

$$\forall e, \forall E \in En^L(e), \exists! \tilde{e} \in E \mid \tilde{e} \in \Sigma' \cap \Sigma^L \tag{20}$$

Because of the above properties, a simple procedure allows to combine $En^r(\cdot)$ and $En^L(\cdot)$ to build $En(e)$ for each $e \in \Sigma' \cup \Sigma^L$.

Given $e \in \Sigma'$, $En(e)$ is simply defined by $En^r(e)$. In fact, if $e \in \Sigma' \setminus \Sigma^L$ clearly $En^r(e) = \emptyset$, while if $e \in \Sigma' \cap \Sigma^L$, $En^r(e) = \{\{e\}\}$ for Property 1.

Instead, for every event $e \in \Sigma^L \setminus \Sigma'$, Property 2 asserts that there is one event $\tilde{e}$ in any $E \in En^L(e)$ which is enabled in the initial state of $A^L$ by the sets of events identified in $En^r(\tilde{e})$. Thus for any $E$ included in some $En^L(e)$ we define $X(E)$ as the “extended $E$”, that is the set of sets built from the combination of all sets in $En^r(\tilde{e})$ and $E$ itself.

$$X(E) = \{E \cup E' \mid \forall E' \in En^r(\tilde{e}), \tilde{e} \in E \cap \Sigma'\} \tag{21}$$

In different words, any set of events $E$ is replaced by a set of sets of events that combines both events in $E$ and all events in $En^r(\tilde{e})$ (that is, the enabling events of $\tilde{e}$).

Finally, we build $En(e)$ for any $e \in \Sigma' \cup \Sigma^L$ as:

$$En(e) = \left\{ \begin{array}{ll} & \{\{e\}\}, e \in \Sigma' \\ & \bigcup_{i=1}^{\lfloor En^L(e) \rfloor} X(E), E \in En^L(e), e \in \Sigma^L \setminus \Sigma' \end{array} \right\}$$ \tag{22}

By applying this procedure to the example in Figs. 4 and 5 we compute sets $En^r(\cdot)$ and $En^L(\cdot)$ for $A'$ and $A^L$ and then combine the obtained results to build $En(y)$. $En^r(\cdot)$ and $En^L(\cdot)$ are the following:

$$En^r(a) = \{\{e\}\} \quad En^L(c) = \{\{e\}\}$$

$$En^r(b) = \{\{a\}\} \quad En^L(d) = \{\{e\}\}$$

$$En^r(c) = \{\{a\}\} \quad En^L(x) = \{\{d\}\}$$

$$En^r(d) = \{\{a\}\} \quad En^L(y) = \{\{c\}, \{d\}\}$$

Then, since $y \in \Sigma^L \setminus \Sigma'$, by (22), we have to compute $X(E_1)$ and $X(E_2)$, where $En^L(y) = \{E_1, E_2\}$ and $E_1 = \{c\}$, $E_2 = \{d\}$. 
For $X(E_1)$, $\check{c} \in E_1 \cap \Sigma$ and $\check{c} = c$. Since $En'(c) = \{\{a\}\}$, then $X(E_1) = \{\{c\} \cup \{a\}\} = \{\{c, a\}\}$. Similarly, for $E_2$ we have $\check{c} \in E_2 \cap \Sigma$, $\check{c} = d$ and $En'(d) = \{\{a\}\}$, then $X(E_2) = \{\{d\} \cup \{a\}\} = \{\{d, a\}\}$. Finally, $En(y) = \{X(E_1) \cup X(E_2)\} = \{\{c, a\}, \{d, a\}\}$

5.2. $En^L$ computation optimization

The divide et impera technique greatly facilitates the $En$ computation, because the enumeration of all paths in both $A'$ and $A^L$ is simpler (the size of both automata is considerably smaller than $A_{sup}$). However, the worst case complexity is still $O(|\Sigma'|!)$ and $O(|\Sigma^L|!)$ for the computation of $En'$ and $En^L$ respectively.

A further improvement in evaluating $En$ can be obtained by observing that, in general, $A' \ll A^L$, or $A'$ is much smaller than $A^L$ in terms of number of states and transitions. In a typical system, in fact, the number of rooms and actions requiring the physical proximity of a user to a device is significantly smaller than the number of actions that can be performed remotely ($|\theta_1| + |\theta_1| \ll |\theta_1| + |\theta_{III}|$). Thus making the enumeration of paths in $A^L$ faster brings significant advantages.

A main consequence of properties introduced in Sect. 5.1 is that, given a target event $e$ and a set of events $E$ enabling two events $a$ and $b$ ($E \in En(a)$ and $E \in En(b)$), $E \cup \{a, b\} \not\in En(e)$ is always true. This means that we can neglect further sequences including both $a$ and $b$, so dramatically reducing the space of states in the automaton exploration.

A solution, which is able to exploit these properties, is shown by Algorithm 1. The algorithm takes the automaton $A^L$ and the target event $e$ as inputs and produces $En^L(e)$ of $\Theta$ as a result. The solution is based on a breadth-first approach where open states are accumulated in a queue $T$. A state $t$ in the queue is defined as a triple $(q, p, \text{cut})$ where $q$ is a state in $Q$, $p$ is the path from $q_0$ to $q$ and $\text{cut}$ is the set of events to be excluded from $\Gamma(q)$.

The exploration of $A^L$ starts from state $(q_0, \varepsilon, \{\})$, where the path to $q_0$ is empty (i.e. $\varepsilon$) as well as the set $\text{cut}$ $(q_0, \varepsilon, \{\})$ is added to $T$ and, as long as the queue is not empty, a state $(q, p, \text{cut})$ is extracted from $T$ and processed.

If the target event $e$ is in $\Gamma(q)$ this means that $pe \in L(A^L)$. In this case, the exploration of the current branch is concluded and $\text{cut}(p)$ is stored in $En^L(e)$. Otherwise, for all events $e' \in \Gamma(q)$ that are not in $\text{cut}$ and are not loops ($\delta(q, e') \neq q$), a new state $t' = (q', pe', \text{cut}')$ is generated, where $q' = \delta(q, e')$ and $\text{cut}'$ includes all events in $\text{cut}$ plus all those enabled in state $q$.

In summary, at each step of the $A^L$ exploration, Algorithm 1 builds a path $p$ and stores a set of events $\text{cut}$ that are no longer considered in the current branch, since they have already been taken into account in some previous step. This set grows with every step, thus reducing significantly the number of events that have still to be evaluated in the following iterations.

The total number of paths depends on the number of new events available at each exploration step. Of course, the best case is when $\Gamma(q_0) = \Sigma \times |\Sigma|$ one event-long paths are then built since, for each state $q' = \delta(q_0, e')$, the set $\text{cut}$ is equal to $\Gamma(q_0)$.

Instead, the worst case occurs when new events be-

### Algorithm 1 Enumeration of $En^L(e)$ for $A^L$

1: Input: $A^L = (Q, \Sigma, \delta, q_0), e$
2: Output: $En^L(e)$
3: $T = \{\}; En^L(e) = \{\}$
4: $t = (q_0, \varepsilon, \{\})$
5: enqueue$(t, T)$
6: while $T \neq \{\} do$
7: $(q, p, \text{cut}) = \text{dequeue}(T)$
8: if $e \in \Gamma(q)$ then
9: $En^L(e) = En^L(e) \cup \mathcal{E}(p)$
10: else
11: for all $e' \in \Gamma(q) \setminus \text{cut} do$
12: if $\delta(q, e') \neq q$ then
13: $q' = \delta(q, e')$
14: $\text{cut}' = \text{cut} \cup \Gamma(q)$
15: $t = (q', pe', \text{cut}')$
16: enqueue$(t', T)$
17: end if
18: end for
19: end if
20: end while
come available and have to be considered at each step of the exploration procedure. In particular, the structure of the automaton with $n$ states is such that: 1. the set $\Sigma$ is partitioned in $\Sigma_1, \Sigma_2, \ldots, \Sigma_{n-1}$ subsets, i.e. $|\Sigma| = \sum_{i=1}^{n-1} |\Sigma_i|$ and 2. for each $e \in \Sigma_i$, $e$ labels a transition from $q_{i-1}$ to $q_i$. In this condition the number of paths that lead to the final state $q_n$ is equal to $\prod_{i=1}^{n-1} |\Sigma_i|$. The maximum value is obtained when the product of all partition sizes is maximum and is approximately equal to $3^{|\Sigma|/3}$. Thus, the complexity of Algorithm 1 in the worst case, is $O(3^{|\Sigma|/3})$.

Algorithm 1 is suitable for exploring $A^L$ and is able to build $En^L(e)$ for each $e \in \Sigma$. It is worth remembering that the same approach cannot be followed to build $En^r(e)$ as properties of $A^L$ do not hold for $A^r$. Fortunately, $A^r$ is much smaller than $A^L$ in practice, so that $En^r(e)$ can be constructed by complete enumeration.

6. Implementation and performance

The methodology introduced in the previous sections has been implemented in a prototype software tool. The Functions computation block in Fig. 2 carries out the operations in Algorithm 1 by means of a purposely developed Java application, able to read the description of $A^L$ and $A^r$ ($A^r||A^L = A_{sup}$) and build the enabling functions $F(a)$. Blocks in the online phase (Fig. 3), have been implemented by two custom Java-based modules, one devoted to refinement and fixing and the other to verification. In particular, the refinement/fixing modules also include the “Z3” off-the-shelf SMT solver \cite{Z3}. In this case, module inputs are the $F(a)$ formulas, that is boolean expressions as needed by the Z3 library.

To provide some preliminary performance evaluation of the proposed approach, some tests were run using a number of synthetically generated scenarios, so as to have a rough estimation of processing times with different sized systems. Test scenarios were generated starting with a small real use case, which is able to stress main aspects of the solution, and by replicating the basic (partial) scheme so as to create several more complex systems of increasing size.

6.1. Basic target system

Fig. 5 shows a number of office areas interconnected through a global network. The number of nodes and high-level policies was progressively incremented by both replicating the basic area and introducing minor adaptations.

Fig. 6 also shows some details about the basic building block structure, that is a simplification of a real use case. $A$ and $B$ are two adjacent rooms and $B$ includes a cabinet $CAB$. Access to the $A$ and $B$ areas is controlled through specific credentials (i.e. physical keys, badges, pin codes, and so on). Room $A$ hosts some general-purpose user workstations ($H_i, i = 1, 2, \ldots, 6$), while a special node ($H_a$) is reserved to administrators for network management activities. Services offered by each $H_i$ are login (accessible either physically or remotely) and backup, while $H_a$ provides an additional (reserved) functionality for remote management of workstations.

Nodes in room $B$ are machines for corporate-level services (i.e., web pages $S_w$, e-mail $S_m$, database $S_{db}$ and code repository $S_c$). Several service functions can be invoked remotely and access to them is guarded by firewall $FW$. Actually, $FW$ is configured so as to enable accesses to $S_{db}$ and $S_c$ from $H_a$ only. Moreover, critical servers are physically confined in cabinet $CAB$. $O$ is the physical environment external to the system, which also includes (part of) the
global interconnection network.

At the beginning any user is assumed to be in $S$, that is outside the system. In this example we consider two different kinds of users: an administrator $adm$ and a generic employee $empl$, characterized by different permissions. Both $adm$ and $empl$ can access $S_w$ and $S_m$, but only the administrator has total control over $S_c$ and $S_d$. In practice, 50 possible actions are defined for each basic office area, and the administrator should be able to perform each one of them, while a smaller set is available to normal users ($S_{empl} \neq \emptyset$).

### 6.2. Performance evaluation

Though the adoption of our methodology relies on all blocks in Figs. 2 and 3 in this paper we focus on the Functions computation and Refinement/fixing elements. Indeed, the experimental evaluation of the time necessary for policy verification showed that, even though it linearly depends on the number of policies, it is almost negligible. As an example, in the worst case, the verification of five hundred policies has proven to take less than 1 ms. The interested reader can find more details about the Automaton generator block and its implementation in [6].

Moreover, when measuring the time needed to carry out the refinement process, performance figures for $empl$ and $adm$ were very similar. Indeed, the complexity of the refinement process performed by the “Z3” solver depends on the number of formulas included in equation [15], which, in turn, corresponds to the number of policies defined for the user (cardinality of $S_w^+$ and $S_w^-$). In our case, it is the same for both users $empl$ and $adm$. For these reasons, and to stress the refinement module capabilities, we reported the performance measures for user $empl$ also varying the actual number of policies ($S_{empl}^+$, $S_{empl}^-$) at each test. The ratio ($|S^+| + |S^-|)/|Actions|$ was changed by defining four policy sets involving 25%, 50%, 75%, and 100% of possible actions respectively.

All tests were performed on a machine equipped with an Intel i7-67000 processor running at 3.40 GHz, 16 GB RAM, and the MS-Windows 10 operating system. Several experiments were carried out to evaluate the Functions computation and Refinement modules, with progressively larger target systems and policy sets, based on the composition of the basic elements discussed above.

A summary of the obtained results is reported in Tab. 1

| Rep. | Rooms | Devices | Actions | C | States | Transitions | Functions (+) | Functions (−) | (|S^+|+|S^-|)/|Actions| |
|------|-------|---------|---------|---|--------|-------------|--------------|---------------|-----------------|
| 1    | 4     | 11      | 53      | 42| 7      | 126         | 0.01         | 0.01           |                 |
| 2    | 7     | 22      | 101     | 83| 12     | 448         | 0.04         | 0.02           |                 |
| 3    | 10    | 33      | 151     | 124| 19     | 1260        | 0.26         | 0.04           |                 |
| 4    | 15    | 44      | 201     | 195| 30     | 3244        | 0.93         | 0.08           |                 |
| 5    | 18    | 55      | 251     | 206| 40     | 7980        | 2.34         | 0.14           |                 |
| 6    | 19    | 66      | 301     | 247| 84     | 18984       | 0.95         | 0.19           |                 |
| 7    | 22    | 77      | 351     | 288| 151    | 44100       | 14.06        | 0.32           |                 |
| 8    | 25    | 88      | 401     | 329| 282    | 160576      | 26.38        | 0.42           |                 |
| 9    | 28    | 99      | 451     | 370| 541    | 220414      | 57.01        | 0.92           |                 |
| 10   | 31    | 110     | 501     | 411| 1056   | 502410      | 94.80        | 0.77           |                 |

Table 1: Performance measurements

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the Refinement procedure is much faster (by two orders of magnitudes) and takes less than one second for the largest test case and the full set of policies.

Note that the Functions computation is carried out only once for the entire system and is not dependent on the number of users. Conversely, the Refinement procedure has to be performed for each user. However, the result of the refinement procedure for one user (i.e., the credential assignment computed for that user) does not affect the refinement procedure for a different user in any way. This means that the processing time for refinement is linearly dependent on the number of users.

6.3. Worst case performance

We also checked the Functions computation module performance in the worst case scenario as discussed in Section 5.2. In particular, several automata were generated with increasing number of actions and a structure able to maximize the number of paths computed for a target event e, which is enabled in the final state of each automaton.

Experimental results are shown in Fig. 9, where the number of actions x varies in the range 5 to 35. The curve \( f(x) = a + b \times 3^{(x \times c)} \), which best fits the experimental data, is also plotted in the figure, with \( a = 9015.639, b = 0.470 \) and \( c = 0.405 \). The measured values are in good agreement with the worst-case asymptotic behavior identified in Section 5.2.

7. Related Works

Several works can be found in the literature concerning policy analysis, refinement and verification. Our contribution mainly focuses on the last two areas.

7.1. Policy Analysis

The main contributions in the area of policy analysis mostly deal with anomaly analysis and policy evaluation.

Anomaly analysis\(^3\) looks for incorrect policy specifications that administrators may introduce in the network. It includes checks for potential errors, conflicts and sub-optimizations affecting either a single policy or a set of security policies\([3]\). Several works have been proposed to

\(^3\)Sometimes anomaly analysis is also referred to as either conflict analysis or policy validation.
perform anomaly analysis for access control policies, that leverage different techniques such as model checking [11], binary decision diagrams (BDDs) [12], graph theory [13], DFSA [14], First Order Logic (FOL) [15] and geometrical model [16].

Policy evaluation, instead, checks whether a request is satisfied/unsatisfied by a set of policies [17]. A typical example of policy evaluation is the change-impact analysis, used to verify the effective impact of inserting, removing and/or modifying a policy. This means comparing original and modified policies. Relevant works on change-impact analysis are [18] and [19], that use BDDs to represent access control policies and perform the analysis for their modifications.

7.2. Policy Verification

Besides systematic literature reviews [20, 21], the number of papers dealing with the verification of access control policies is not large, especially when it is compared to the amount of literature dealing with policy analysis. Notable works in this area are [22, 23, 24, 25, 26].

In [22, 23, 24], authors propose the automatic verification of access control policies against a set of properties. Verification is carried out by translating properties into a boolean satisfiability problem to be managed by means of a SAT solver. Differently from our approach, [22] and [23] model the target system in terms of access control policies and check if properties are violated by policies. Our solution, instead, builds on a fine-grained model of the real system where policies have to be verified. The SAT solver is used for refinement and fixing. In [26] authors propose a complete methodology based on BDD for the specification, verification and enforcement of access control policies. Differently from our approach, authors do not propose any methodology to fix the identified anomalies. Other tools such as NP-View and PolicyGlobe, respectively described in [24] and [25], allow for checking global policy implementations by processing information about the network topology and configuration. Unfortunately, this is limited to systems where policy enforcement mechanisms are assumed to be available. This reduces the tools usefulness as, for instance, they can produce meaningful results only for systems consisting of nodes running SE-Linux.

Note that, at the moment our approach does not allow to model and check the correct implementation of policies that require counting the number of times a certain action is performed by one or multiple users (similar, e.g., to history-based policies [27, 28, 29]). Indeed, our approach is meant to be used for a priori system verification/configuration and not as a runtime tool for access control. As such, while the model is able to capture action dependencies determined by system characteristics and configuration, it does not describe either multiple user interactions nor the flowing of time.

7.3. Policy Refinement

Although refinement is one of most ambitious goals in policy-based management when access control is considered, relatively few works have been published in the last decade, with particular focus in this area. Refinement of access control policies was originally introduced in [30, 31, 32], but little or no development has appeared since then.

More recently, with the advent of both the software defined networking (SDN) and network function virtualization (NFV) paradigms, interest in policy refinement techniques has revived [33, 34, 35]. It is worth noting that while other proposals aim at defining high-level specification languages, which are able to describe access control policies and translate them into low-level configuration commands, our solution exploits refinement techniques to search all sets of credentials that satisfy the policies automatically. If no solution is found, changes in either the policy specification or the system structure become the only viable alternatives.
8. Conclusions and Future work

In this paper, a comprehensive approach has been presented to solve some common problems faced by network administrators when dealing with access control through policy-based management. The proposed approach can be exploited to:

1. find a user credential assignment that correctly enforces a set of access control policies (policy refinement);
2. complete a partial user credential assignment so that results correctly enforce the specified policies (constrained policy refinement);
3. check whether a specific user credential assignment makes the system behavior match the policy specification (policy verification). If this is not the case, the approach can be used to introduce automatic changes to fix the highlighted discrepancies between the actual and desired behavior (i.e. anomaly fixing).

Our technique is based on the definition of models for both the access control policies and the target networking system (in terms of topology, s/w and h/w resources and how they are accessed by users). The system model is then processed to deduce what sequences of actions are needed to perform a specific operation. The resulting information is used to determine the minimal sets of credentials (if any) that shall be owned by users to perform all and only the actions assigned to them by the policies.

The feasibility of the approach has been verified by means of a prototype s/w tool, which is based on Java and the Z3 library, an off-the-shelf SMT solver. The tool performance was tested in a number of realistic network scenarios and its scalability also checked (the tool is able to provide solutions in a very short time even for networks consisting of several hundred nodes and target resources).

In the future, we plan to extend the expressiveness of the model to take into account further types of network devices, such as VPN gateways and intrusion detection systems (IDSs). Furthermore, we intend to apply this methodology so as to consider not only user credential assignments but also other aspects of network configuration (e.g., firewall rules). Finally, it is our intention to perform some empirical assessment to evaluate the ability of the model to capture the nature of real networks and the effectiveness of the approach in helping administrators managing access control in their networks.

References


Manuel Cheminod received the M.S. in 2005 and the Ph.D in Computer Engineering in 2010, both from the Politecnico di Torino, Italy. He has a Researcher with the Institute of Electronics, Computer, and Telecommunication Engineering of the National Research Council of Italy, in Torino. His research interests include formal methods for the verification of security properties in industrial systems and the exploitation of SDN/NFV paradigms to improve the security level of the same critical systems.

Luca Durante is Senior Researcher at the Institute of Electronics, Information Engineering and Telecommunication of the National Research Council of Italy (CNR-IEIIT) in Torino, Italy. He has co-authored about 60 scientific journal, conference papers and technical reports in the area of industrial communication protocols and formal techniques for distributed systems. He also serves as a technical referee for several international conferences and journals, and he has been the scientific coordinator of the CNR team in European Projects. He served as Program Co-Chairman for the 13th (WFCS 2017) edition of the IEEE Workshop on Factory Communication Systems. Currently his research interests include formal verification of cryptographic protocols, source-level model checking of software and network vulnerability and dependability analysis.

Lucia Seno (M’15) received the B.S. and M.S. degrees in Automation Engineering and the Ph.D. degree in Information Engineering from the University of Padova, Padova, Italy, in 2004, 2007, and 2011, respectively.

Since then, she has been with the Institute of Electronics, Computer, and Telecommunication Engineering of the National Research Council of Italy in Padova, and, subsequently, in Torino, Italy, where she is currently a Researcher. Her research interests include industrial communication systems and technologies, wireless networks and wireless sensor networks for real-time and safety-critical control applications, formal verification of system security and communication protocols, and model-checking.

Fulvio Valenza received the M.Sc. (summa cum laude) in 2013 and the Ph.D. in Computer Engineering in 2017 from the Politecnico di Torino, Torino, Italy. His research activity focus on network security policies. Currently he is a Researcher at the CNR-IEIIT Torino, Italy, where he works on orchestration and management of network security functions in the context of SDN/NFV-based networks and industrial systems.

Adriano Valenzano (SM’09) received the Laurea degree in electronic engineering from Politecnico di Torino, Torino, Italy, in 1980. He is Director of Research with the National Research Council of Italy (CNR). He is currently with the Institute of Electronics, Computer and Telecomm-
munication Engineering (IEIIT), Torino, Italy, where he is responsible for research concerning distributed computer systems, local area networks, and communication protocols. He has coauthored approximately 200 refereed journal and conference papers in the area of computer engineering. Dr. Valenzano is the recipient of the 2013 IEEE IES and ABB Lifetime Contribution to Factory Automation Award. He also received, as a coauthor, the Best Paper Award presented at the 5th, 8th and 13th IEEE Workshops on Factory Communication Systems (WFCS 2004, WFCS 2010 and WFCS 2017). He has served as a technical referee for several international journals and conferences, also taking part in the program committees of international events of primary importance. Since 2007, he has been serving as an Associate Editor for the IEEE Transactions on Industrial Informatics.