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Original

Trusted GNSS-Based Time Synchronization for Industry 4.0 Applications / Margaria, Davide; Vesco, Andrea. - In: APPLIED SCIENCES. - ISSN 2076-3417. - ELETTRONICO. - 11:18(2021), p. 8288. [10.3390/app11188288]

Availability:

This version is available at: 11583/2922136 since: 2021-09-08T10:54:16Z

Publisher:

MDPI, Basel, Switzerland

Published

DOI:10.3390/app11188288

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Trusted GNSS-based Time Synchronization for Industry 4.0 Applications

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Abstract: The protection of satellite-derived timing information is becoming a fundamental requirement in Industry 4.0 applications, as well as in a growing number of critical infrastructures. All the industrial systems where several nodes or devices communicate and/or coordinate their functionalities by means of a communication network need accurate, reliable and trusted time synchronization. For instance, the correct operation of automation and control systems, measurement and automatic test systems, power generation, transmission, and distribution typically require a sub-microsecond time accuracy. This paper analyses the main attack vectors and stresses the need for software integrity control at network nodes of Industry 4.0 applications to complement existing security solutions that focus on GNSS RF Spectrum and Precise Time Protocol (PTP), also known as IEEE-1588. A real implementation of a Software Integrity Architecture in accordance with Trusted Computing principles concludes the work together with the presentation of promising results obtained with a flexible and reconfigurable testbed for hands-on activities.

Keywords: trusted computing; industry 4.0; cyber-physical system; time synchronization; cybersecurity; global navigation satellite system; embedded system

Citation: Margaria, D.; Vesco, A. Trusted GNSS-based Time Synchronization for Industry 4.0 Applications. *Appl. Sci.* **2021**, *11*, 0. <https://doi.org/>

Academic Editor: Silvio Abrate

Received:

Accepted:

Published:

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1. Introduction

The cornerstone of Industry 4.0 [1] are Cyber-Physical Systems (CPS) which are closely connected with computer systems and which can interact and collaborate with other CPS systems [2]. These concepts represent the basis of decentralization and cooperation between connected systems: they are closely linked to the Industry 4.0 paradigm and fully enabled by Information and Communication Technologies [3–5].

Networked environments and Industrial Internet of Things (IIoT) [4] need accurate, reliable and trusted time synchronization and time distribution [6], with different requirements in terms of synchronization accuracy and precision. For instance, the scheduling and synchronization of multiple factories represents an interesting high level application of the Industry 4.0 paradigm [5]. Due to rapidly changing market requirements, factories have shifted from a centralized to a more decentralized structure in many areas of decision making, including scheduling. Since limited resources make scheduling an important decision in the production, accurate time synchronization and efficient scheduling solutions are vital for improving the productivity in a multi-factory production network [7].

Another relevant example is the case of Digital Twins (DTs) in Industry 4.0 [8]: the underlying idea is that a real product and its virtual counterpart are twins that travel a parallel journey from design and development to production and service life. A DT must be always in sync with the corresponding physical asset. Such synchronization is needed to correctly update and to keep the operational data (*i.e.*, failure and erroneous data) consistent within the production system.

37 Nowadays, several Industry 4.0 applications already imply stringent synchroniza-
 38 tion requirements (e.g., sub-microsecond accuracy). Among the others, we consider the
 39 following ones as especially relevant:

- 40 • *Automation and control systems* [9], that need an accurate common notion of time
 41 as well as a reliable shared communication medium for timely data exchange,
 42 for instance to synchronize multi axis drive systems and subsystems with cyclic
 43 operation;
- 44 • *Measurement and automatic test systems* [10], that usually take advantage of accu-
 45 rate time stamping of logged data, for example to correlate acquired values in
 46 decentralized locations;
- 47 • *Power generation, transmission and distribution systems* [11,12], requiring a precise
 48 time synchronization of all the critical points within the power grid to accurately
 49 measure the delivered/consumed power and to predict critical load situations.

50 In this context, satellite-derived timing information plays a key role in the pro-
 51 visioning of an absolute time reference to a significant number of current and future
 52 connected systems. *Global Navigation Satellite System* (GNSS) receivers combined with
 53 specific transport protocols or with distributed synchronization approaches can indeed
 54 satisfy the stringent requirements foreseen in several industrial applications [13].

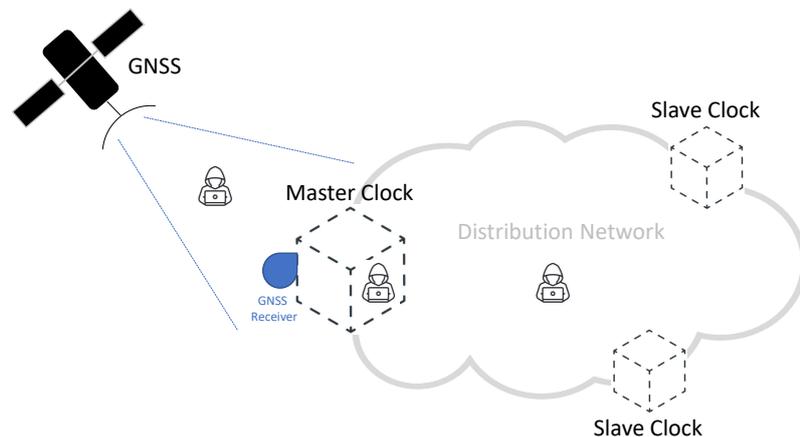


Figure 1. Simplified scheme of a GNSS-based synchronization distribution network.

55 Figure 1 provides a common scheme of a synchronization network, representative
 56 of a typical time distribution network and highlighting the role of GNSS technologies.
 57 In this scheme, accurate time information is generated by each GNSS satellite and then
 58 transmitted by means of properly formatted Radio-Frequency (RF) signals [14,15]. A
 59 Master Clock node of the network accurately synchronizes its local clock using a GNSS
 60 receiver, taking advantage of the signals received from the GNSS satellites visible at its
 61 location [16]. As illustrated in Figure 1, a distribution network allows to distribute the
 62 timing information to all the nodes, achieving an accurate synchronization between the
 63 Master Clock and multiple Slave Clocks. The distribution network can take advantage
 64 of different protocols and technologies, depending on the application requirements
 65 and constraints. Among the others, the *Precision Time Protocol* (PTP) represents a well-
 66 known and widely adopted solution for accurate time distribution in a network of
 67 clocks organized in a master-slave hierarchy [16,17]. Such synchronization protocol, also
 68 known as IEEE-1588 standard [18,19], allows for absolute time synchronization in the
 69 range of hundreds of nanoseconds through hardware assistance (SyncE) and, potentially,
 70 sub-nanosecond accuracy with the White Rabbit extension of PTP (WR-PTP) [20,21].

71 Nevertheless, a conscious adoption of the Industry 4.0 paradigm also requires the
 72 introduction of appropriate cybersecurity solutions. In fact, cyberattacks and hacking
 73 of factory industrial control systems are dramatically increasing in recent years. Proper
 74 mitigation measures against these attacks are needed to avoid the emergence of an

75 internet of *insecure* industrial things [22]. The same statement is also applicable to all
76 the grids and telecommunications networks recognized as critical infrastructures [23,24],
77 especially the upcoming 5G [16,17,25] that is considered as a key enabler for several
78 Industry 4.0 applications. In all these cases, possible synchronization inaccuracies can
79 directly result in a degradation of the quality of service provided by the communication
80 network (*e.g.*, reduced throughput, increased latency and jitter) and, in extreme cases, in
81 a complete disruption of the considered system or application.

82 Two areas in Figure 1 represent likely targets for potential attacks (*i.e.*, viable *attack*
83 *vectors*) and, thus, need appropriate protection:

- 84 1. the **GNSS RF spectrum**, potentially targeted by attacks against the GNSS signals;
- 85 2. the **Time Distribution Network**, potentially targeted by attacks to the synchroniza-
86 tion information. We can further divide this category into:
 - 87 (a) attacks targeting the Precision Time Protocol over the network, and
 - 88 (b) attacks against the integrity of PTP software (SW) running at each node and
89 the related configuration files.

90 It is worth noting that recent scientific literature has extensively investigated attacks
91 to the GNSS RF spectrum (*i.e.*, attack vector 1) and to the PTP protocol and packets
92 over the network (*i.e.*, 2.a). Proper countermeasures against them are already available
93 (*e.g.*, refer to [26–28] and references therein). For this reason, following sections will
94 summarize these specific attack vectors, while their experimental assessment is out of
95 the scope for this paper.

96 On the other hand, to the best of the authors' knowledge, previous works have not
97 widely covered the attacks to the PTP SW integrity within nodes (*i.e.*, 2.b). This category
98 of attacks is gaining special attention in both scientific and industrial communities,
99 where the investigation of possible cybersecurity solutions is an active research topic,
100 stressing the need for software integrity control at network nodes. For these reasons,
101 we recognize this category of attacks and related solutions as especially relevant and
102 challenging for Industry 4.0 applications and, thus, worth of further investigation.

103 Motivated by this background, the paper has the following objectives:

- 104 • an analysis of attack vectors for GNSS-based synchronization distribution networks,
105 with a special attention to attacks not widely covered in prior work (*i.e.*, 2.b);
- 106 • the proposition of a novel solution for PTP nodes augmented with *Trusted Computing*
107 technologies to counteract attacks to their integrity, complementing existing security
108 solutions that focus on GNSS RF Spectrum and PTP protocol only;
- 109 • a reference implementation on a testbed, capable to demonstrate the proposed
110 solution by means of quantitative results.

111 Our proposed approach differs from other *authentication and authorization-based*
112 solutions already available in prior work and suitable to ensure the integrity of the
113 network protocols and all type of exchanged packets. In fact, our solution leverages
114 the *Trusted Computing* principles to achieve a level of trust in the behaviour of the
115 synchronization distribution network. It is capable to ensure the integrity of the soft-
116 ware and configuration of all the nodes and, thus, to avoid the exchange of incorrect
117 synchronization information over secure protocols. For a more complete view on prior
118 works about SW vulnerabilities in other domains, interested readers can also refer to
119 [29–31] and the references therein.

120 In this sense, the contribution of the paper is twofold:

- 121 • We present our proposed implementation based on *Trusted Computing* and capable
122 of protecting the integrity of the nodes of a synchronization distribution network.
- 123 • Next, we propose and describe a new testbed, suitable to emulate the identified
124 attacks and to demonstrate the effectiveness of the proposed solution.

125 After this introduction, the paper is organized as follows. Section 2 provides an
126 overview of the vulnerabilities and solutions related to the GNSS RF spectrum, while
127 Section 3 covers the network attack vector. After that, Section 4 discusses the *Trusted*

128 Computing paradigm, as a suitable solution to protect the SW integrity of the nodes.
129 Then, Section 5 describes in detail our implementation and Section 6 presents and
130 discusses the obtained results. Finally, Section 7 concludes the paper with the main
131 remarks.

132 2. GNSS RF Spectrum Attack Vector

133 GNSS technologies [14,15] have shown a remarkable growth in the last years, being
134 nowadays adopted in the most different fields of applications such as: consumer devices
135 (*e.g.*, smartphones, cameras, wearable devices, etc.), vehicular applications (*e.g.*, road user
136 charging, bike sharing, connected and automated driving, etc.), manned and unmanned
137 aviation (*e.g.*, drones), industrial applications and even in critical infrastructures. Every
138 system that makes use of GNSS-derived data must trust them before taking decisions,
139 especially in case of safety-critical or liability-critical applications.

140 The widespread adoption of GNSS technologies creates incentives for the attackers
141 that want to impair or fool any system that rely on GNSS to estimate the position, the
142 velocity and, especially, the *time* information [26]. GNSS receivers and connected devices
143 integrating these receivers are all vulnerable to intentional attacks, aiming to affect the
144 availability and the reliability of the GNSS signals and data. In general, GNSS RF attacks
145 are classified in three main categories [24,26,27]:

- 146 1. **Jamming**, that is the blocking of the reception of GNSS signals by intentionally
147 emitting RF interferences to disrupt the functionalities of the receivers, in order to
148 reduce the signal-to-noise power level;
- 149 2. **Meaconing**, that corresponds to the rebroadcasting of delayed GNSS signals, with-
150 out any distinction between signals received from different satellites;
- 151 3. **Spoofing**, that refers to the transmission of counterfeit GNSS-like signals, with the
152 intent to produce false position and/or time data at the victim receiver.

153 Each of these three categories can be put in place in different ways, implying a
154 different complexity and cost at the attacker side. Among the possible attacks, only few
155 can be considered likely to be deployed in real applications [12,16,24].

156 Different countermeasures against these attacks are already available and widely
157 discussed by the GNSS research community [26]. Please note that a comprehensive
158 review of the state of the art related to GNSS RF attacks and proposed solutions is out of
159 the scope for this paper. Nonetheless, interested readers can refer to [17,26,27] and the
160 references therein.

161 Among the most recent solutions, it is worth to point out the on-going effort of
162 the European Galileo program to gradually add authentication services to its first and
163 second generation of satellites signals, in order to enable authentication functionalities
164 for future civil receivers [27]. The Galileo *Open Service Navigation Message Authentication*
165 (OSNMA) consists in a mechanism that allows the receivers to verify the authenticity
166 of GNSS information, making sure that the data they receive are indeed from Galileo
167 satellites and have not been modified in any way [32]. GNSS receivers can take ad-
168 vantage of such capability for implementing simple but effective detection methods
169 against several spoofing attacks [16,24,33]. The first-ever signal-in-space transmission of
170 Galileo OSNMA started on November 2020 and tests are currently underway in order to
171 consolidate the service [34].

172 3. Time Distribution Network Attack Vector

173 The time distribution network represents a potential target for cyberattacks that can
174 have an interest on degrading or disrupting the availability, integrity and reliability of
175 the exchanged timing information.

176 In fact, several security mechanisms typically have a critical interdependence with
177 timing synchronization. For instance, security solutions based on signed keys or certifi-
178 cates require accurate timing to determine whether they are valid and for how long. In
179 this sense, accurate time synchronization is needed to establish a valid security system

180 and, on the other hand, a valid security system is required to confirm the accuracy of the
181 timestamps [28]. The following paragraphs discuss the network attack vector, focusing
182 on two sub-categories: (i) attacks against the PTP protocol used on the distribution
183 network, and (ii) attacks against the PTP SW integrity and PTP configuration files of
184 each node of the network.

185 3.1. Attacks Against the PTP Protocol and Packets

186 Recent scientific literature report several documented examples of attacks against
187 timing protocols, including attacks against PTP [28,35,36]. These attacks are usually
188 categorized as Denial-of-Service (DoS) attacks, feasible at various network layers, and
189 Man-In-The-Middle (MITM), including clock masquerade, replay and filtering attacks
190 [28,35,36]. For instance, Alghamdi and Schukat [37] proposed specific attack strategies
191 and implementations: packet content manipulation, packet removal, packet delay ma-
192 nipulation, time source degradation, master spoofing, slave spoofing, compromised
193 Best Master Clock Algorithm (BMCA), packet replay, and DoS. In addition, DeCusatis
194 *et al.* [28] recently proposed and experimentally demonstrated a novel class of insider
195 threats, including two variants of DoS spamming attacks, capable to incorrectly steer or
196 permanently skew the slave's clock, and a master clock takeover attack.

197 Mitigation techniques and solutions against the previous categories of attacks have
198 also been proposed and discussed. It is worth to mention that IEEE-1588-2008 standard
199 [18] defined an experimental security extension (Annex K) in order to protect a PTP
200 network. However, a number of weaknesses and drawbacks in this approach have been
201 identified and, today, Annex K is deprecated [35]. Further on, a new version of the stan-
202 dard IEEE-1588-2019 has recently been published. It contains a new security extension in
203 the Annex P, based on a multi-pronged approach [19]. Apart this remarkable standard-
204 ization effort, the scientific literature includes other recent solutions. Among the others,
205 it is worth to point out a proposed extension for PTP of the key management mechanism
206 used in Network Time Security (NTS) [38] and an identity-based authentication system,
207 *i.e.*, First Packet Authentication with Transport Access Control [28].

208 3.2. Attacks Against the SW Integrity and Configuration of the Nodes

209 The protection of the timing information over the GNSS RF spectrum and in transit
210 over a PTP network against cyberattacks is not enough to ensure trust and rely on the
211 GNSS precise timing. In fact, even assuming a state-of-the-art distribution network as
212 in Figure 1, including all the previously cited security solutions, such system would
213 be still vulnerable to potential attacks against the software integrity of Master and/or
214 Slave nodes. Specific attacks targeting the integrity of the SW stack running within those
215 nodes must also be considered to avoid that nodes exchange *incorrect synchronization*
216 *information over secure protocols*.

217 Figure 2 illustrates a possible SW stack running on Master Clock and Slave Clock
218 nodes. We adopt the simplified scheme in Figure 2 as a study case, intended to be
219 as general as possible and representative of the typical SW modules installed and
220 configured on real PTP nodes, but without the ambition to cover all the possible SW
221 configurations and different hardware components (*e.g.*, Network Interface Card with
222 or without PTP hardware timestamping, internal clock steering done via SW or with a
223 dedicated FPGA, etc.).

224 We propose two possible configurations for the Master node on the left side and
225 on the central part of Figure 2, respectively. Both these configurations include a GNSS
226 module capable to receive the RF signals from one or multiple satellite constellations
227 (*e.g.*, GPS, Galileo, GLONASS, and/or BeiDou) in order to estimate Position, Velocity,
228 and Time (PVT) information. The GNSS receiver typically provides two outputs:

- 229 1. a textual interface over a serial communication protocol, providing the PVT data
230 coded as "sentences" according to the National Marine Electronics Association
231 (NMEA) 0183 standard [39], and

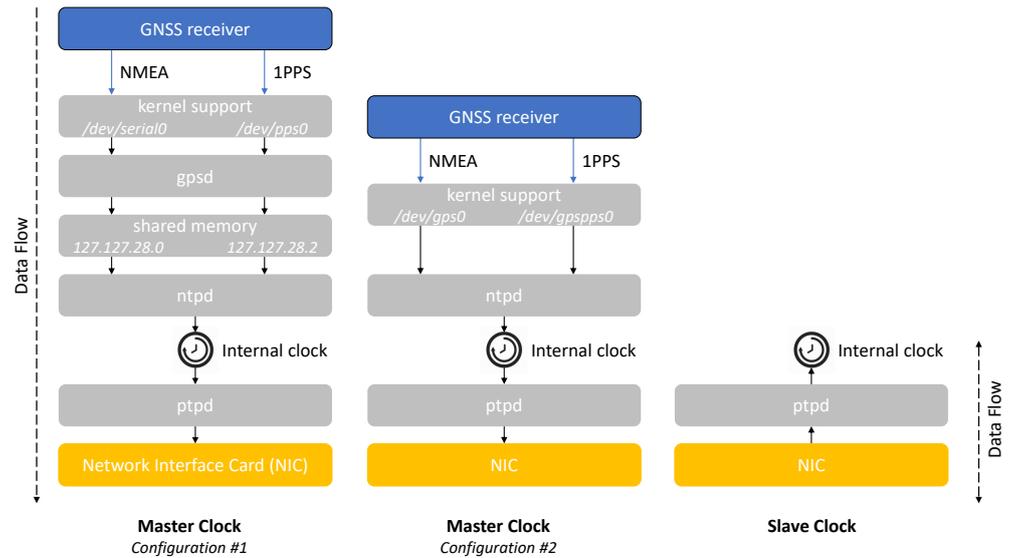


Figure 2. Possible SW stack and configurations for Master Clock and Slave Clock nodes.

232 2. the 1 Pulse-Per-Second (1PPS), that is a high precision analog signal having leading
 233 pulse edges synchronous with the beginning of each second of the time scale.

234 Different SW tools and/or daemons can accept as inputs the NMEA and 1PPS, taking
 235 advantage of the Linux kernel support to make them visible as devices [40].

236 In the configuration #1, these devices are a serial port for NMEA (/dev/serial0)
 237 and the Linux Pulse Per Second Application Programming Interface (PPSAPI) for 1PPS
 238 (/dev/ppp0) [40]. A service daemon capable to parse the GNSS data (*i.e.*, gpsd [41]) uses
 239 both of them as inputs. We also shared such inputs to ntpd, a daemon implementing the
 240 Network Time Protocol (NTP) version 4 [42], by means of two shared memory segments
 241 (*i.e.*, 127.127.28.0 and 127.127.28.2) properly configured using the Shared Memory
 242 Driver of ntpd [43]. In this way, ntpd is capable to correctly discipline the internal clock
 243 of the Master node, ensuring its accurate synchronization with respect to the GNSS time
 244 scale.

245 The configuration #2 represents a simplified alternative without gpsd. In this case,
 246 we directly provide the NMEA and 1PPS outputs of the GNSS receiver to ntpd by
 247 means of two symbolic links to the actual devices (*i.e.*, /dev/gps0 and /dev/gpspps0,
 248 respectively). It is worth to notice that ntpd can use these symbolic links as two separate
 249 inputs, taking advantage of two dedicated drivers (*i.e.*, the Generic NMEA GPS Receiver
 250 [44] and the PPS Clock Discipline [45]) or as a single NMEA stream directly disciplined
 251 by the 1PPS (*i.e.*, by properly setting the fudge flag1 of the Generic NMEA driver
 252 to enable the PPS signal processing [44]). The second option is preferable in order to
 253 simplify the configuration of ntpd, thus we will use it in the following discussion.

254 Once ntpd has correctly synchronized the internal clock of the Master node, both
 255 the configurations #1 and #2 adopt the PTP daemon (*i.e.*, ptpd [46]) to distribute the time
 256 synchronization over the network. In accordance to PTP working principles, ptpd can
 257 accurately synchronize multiple Slave nodes to the internal clock of the Master node.
 258 In detail, all the Slave nodes reach the PTP_SLAVE status [18,19], acquire the control and
 259 start correcting their internal clocks. It is known that ptpd can achieve microsecond level
 260 of time coordination, even on limited platforms [46].

261 The protection of the integrity of the SW stack in Figure 2 is of paramount impor-
 262 tance to avoid that nodes share intentionally modified timing information over any
 263 secure version of PTP. Aiming to detect possible intentional changes in the work logic of
 264 the Master and/or the Slave node, next section proposes a software integrity architecture
 265 based on Trusted Computing principles and customized for embedded devices.

266 4. Trusted Computing

267 The Trusted Computing paradigm can be implemented in different ways and as a
 268 combination of different approaches and trust decisions to build the Secure Boot and/or
 269 Remote Attestation. The following paragraphs present our implementation for the
 270 purpose of protecting the integrity of an embedded system operating as a node of a
 271 synchronization distribution network.

272 4.1. Trusted Computing Base

273 Trusted Computing (TC) is used in this paper as per the Trusted Computing Group's
 274 (TCG) definition: a set of interoperable technologies to achieve a level of trust in the
 275 behaviour of an embedded system. The core element of TC is the hardware Root of
 276 Trust called Trusted Platform Module (TPM). A TPM is a tamper resistant piece of
 277 cryptographic hardware integrated with the system board that implements primitive
 278 cryptographic functions on top of which more complex features can be built in accor-
 279 dance with TCG specification TPM 2.0 [47]. The final objective for the application of TC
 280 in this work is to enable nodes to measure and prove cryptographically their integrity, *i.e.*,
 281 prove that the software running is the intended one and it has not been tampered with,
 282 to the other nodes involved in the synchronization distribution. Integrity measurement
 283 is therefore the process the nodes adopt to collect and digest the information about the
 284 integrity of their software and configuration for the purpose of being attested/verified.
 285 The three main hardware Roots of Trust to make the node a Trusted Computing Base
 286 (TCB) are:

- 287 1. Root of Trust for Measurements (RTM): the Core Root of Trust for Measurements
 288 (CRTM) that act as the Trust Anchor; it is commonly implemented by the initial
 289 bootloader secure code executed at power up or hardware reset;
- 290 2. Root of Trust for Storage (RTS): the TPM that provides (i) a set of Platform Configu-
 291 ration Registers (PCRs) to securely accumulate the integrity measurements in form
 292 of hashes, (ii) a key hierarchy architecture to securely store objects protected via
 293 encryption by the TPM on the file system;
- 294 3. Root of Trust for Reporting (RTR): the TPM that signs the values of a PCR set
 295 using an Attestation Key bound to the Endorsement Key, *i.e.*, a resident key that
 296 represents the TPM identity and that guarantees the origin and the integrity of the
 297 PCR values shared for the purpose of attestation.

298 4.2. Chain of Trust

299 The identification of the software components is performed through a Chain of
 300 Trust, *i.e.*, a hierarchy of trust rooted into a Trust Anchor (TA). The TA is an authoritative
 301 entity for which trust is assumed and not derived, *i.e.*, usually a small and carefully
 302 designed piece of firmware, executed by the system at power up or hardware reset. The
 303 TA is responsible for starting to build the Chain of Trust during bootstrap of the node.
 304 Each element of the chain is a software component that when executed performs its
 305 tasks and then identifies, loads and finally executes the next software component (see
 306 Figure 3).

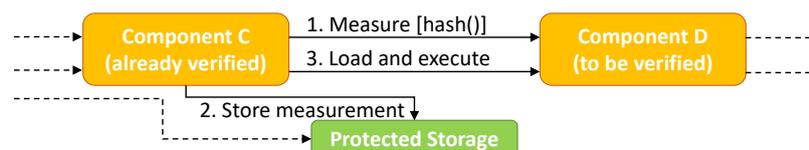


Figure 3. Basic Element of a Chain of Trust.

307 The identification of a component is made of two steps: the measurement, *i.e.*, a
 308 digest, calculated over that component and the storage of such measurement in the TPM,
 309 *i.e.*, the accumulation of the measurement via PCR extension operation.

Each PCR can contain one digest value which is the result of one or more PCR extension operations: $PCR_{i+1} = PCR_i || hash(component)$. For instance, the hash function can be sha1, sha256, or sha384. The Infineon TPM used in our prototype [48] implements sha1 and sha256 thus, for the sake of simplicity, we selected sha1. It is not possible to store a value directly in a PCR, but the only way to access a PCR for writing is through the PCR extend operation: this guarantees that all measurements are accumulated. By building on the fact that the Trust Anchor is trusted by definition and each software component is considered trusted at least to identify the next component, the identification of all components is trusted as well due to the transitive property of trust. The bootstrap process enhanced to build the Chain of Trust, *i.e.*, the capability of each software component to identify and load the next one, is called Authenticated Boot(strap). The Authenticated Boot goes from the CRTM up to the applications in userspace through bootloader and Linux kernel.

4.3. Trust Decision

The Trust Decision is about deciding whether a platform can be considered trusted for an intended purpose or not. It builds upon the Authenticated Boot, can occur (at any meaningful time) during it or after it and implies the cryptographic verification of the loaded software components to identify them and verify that they are the expected ones. If the Trust Decision is taken during the bootstrap, the verification is performed by the platform itself (usually using the TPM) over the components loaded until the decision time: based on the result of the verification, the bootstrap process can be stopped. This portion of the Authenticated Boot(strap) is called Secure Boot(strap). If the verification is performed after the bootstrap completion, it is called Remote Attestation, as it requires a remote entity, the Attestor, that performs the verification with the support of the trusted component on the platform, *i.e.*, the TPM. In our prototype we implemented both the Secure Boot and the Remote Attestation.

4.4. Secure Boot

The Secure Boot portion can be implemented by means of a combination of the (full) disk encryption with the TPM sealing of the encryption/decryption key to a specific set and configuration of the components loaded and executed before mounting the encrypted partition. Therefore, the availability of the encryption/decryption key is bound to specific values of a set of PCRs containing the accumulated measurements of the loaded components. If such components (and/or their configurations) are different from the expected ones, this is reflected into the PCR values different from those stored along with the encryption/decryption key in a Non-Volatile (NV) storage index. Indeed it is possible to associate a set of PCRs and their values representing a specific configuration to the disk encryption/decryption key when the data is protected (sealing). The complementary operation is the unsealing when the TPM checks whether the current PCR values meet the values' set stored along with the disk encryption/decryption key. If they differ, the TPM does not release the encryption/decryption key (it does the Trust Decision) and, since the encrypted partition cannot be mounted, the bootstrap process gets stopped, thus implementing the Secure Boot. In our prototype we implemented the Secure Boot using a modified version of Cryptsetup [49] that uses a NV storage index to protect the encryption/decryption phase.

4.5. Integrity Measurement Architecture (IMA)

Integrity Measurement Architecture (IMA) [50,51] is a security subsystem of the Linux kernel. When enabled, it measures all files that match a given policy (see Appendix A.1) that are loaded and possibly executed. All measurement records are kept in memory in the so-called IMA log and an overall integrity checksum is retained using the TPM PCR10. This is an excerpt of an IMA log for the Linux kernel v.4.19 we used:

```
PCR template-hash          tpl filedata-hash          file-pathname
10 03eb317e687cbd21e360e257b4810f8ac711fb4c ima 9797edf8d0eed36b1cf92547816051c8af4e45ee boot_aggregate
10 1110984f14fc70c87f90053612c3feaa068f66d6 ima 339c9d9d10d1fc25731d6f3d00b80e173e35456c /lib/systemd/systemd
10 bc447ab6e2f476455644dbcf9187c922418f02ba ima 369c4027e9b09131c6971ee963bfd37d41dc251c /lib/arm-linux-gnueabi/hf/ld-2.28.so
10 e359bd4b3a5b64f125dda645958237a123fe9fe7 ima 9e7916b2b9232f7db4cff863a6588e10253c0285 /etc/ld.so.preload
...
```

360 Each measurement record includes:
 361 IMA-record = PCR-index || template-hash || tpl || filedata-hash || file-pathname,
 362 where:
 363 • PCR-index is the index of the used PCR, with default value equal to 10, *i.e.*, the PCR
 364 extended, one-by-one, by template-hash;
 365 • tpl is the template name;
 366 • template-hash = sha1(filedata-hash length, filedata-hash, ...
 367 file-pathname length, file-pathname);
 368 • filedata-hash = sha1(filedata).
 369 IMA also includes the function *Appraisal* that together with the *Linux Extended*
 370 *Validation Module* (EVM) implements a fine grained Secure Boot process. In our prototype,
 371 to implement the Secure Boot we selected the approach based on an encrypted partition
 372 and the encryption/decryption key sealed to the TPM and a set of PCR values.

373 4.6. Remote Attestation

374 The Trust Decision can also occur when the authenticated bootstrap is completed;
 375 with this option a remote node, *i.e.*, the *Attestor*, is responsible to challenging and
 376 verifying the trustworthiness of target nodes with a given periodicity, as depicted in
 Figure 4. The Attestor performs a Remote Attestation by asking the TCB to be attested

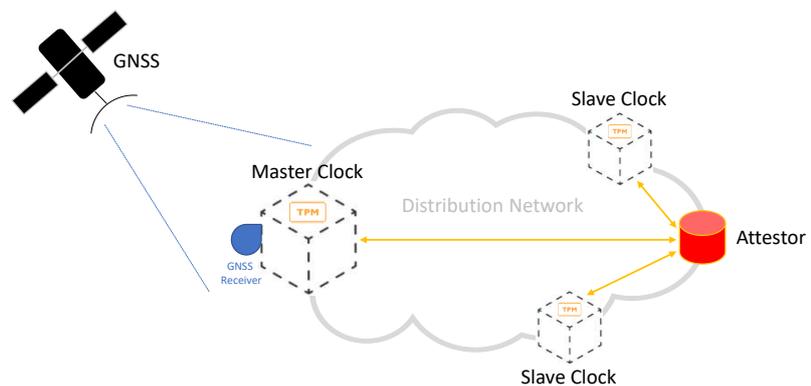


Figure 4. Illustration of Remote Attestation on the synchronization distribution network.

377 to send information about the loaded components for the cryptographic verification.

378 In this work the Remote Attestation is implemented in its easiest form. The Attestor
 379 opens a standard secure channel with mutual authentication by means of TLS [52] with
 380 the target TCB to protect the communication. Then, the Attestor challenges the TCB
 381 and verifies its trustworthiness. To avoid attacks where a malicious node can fool the
 382 Attestor by acting as MITM and acquire the attestation data from a trusted node, the
 383 node implements a soft binding between the TPM identity and the Secure Channel
 384 identity. This is implemented by binding the Attestation Key (AK) used to perform the
 385 *TPM_Quote* operation over the selected PCRs values and the key used for authentication
 386 during secure channel handshake. More practically, this is done by using the Public
 387 Key Certificate of the key used for authentication to extend one PCR of the TPM. The
 388 *TPM_Quote* is substantially a digital signature over a set of relevant PCR values and the
 389 nonce sent by the Attestor for freshness purposes to avoid replay attacks.

390 Once the TLS channel is established between the two parties that successfully identified
 391 each other, the Remote Attestation takes place over it according to the following
 392 sequence of actions:
 393

- 394 1. The Attestor sends the nonce to the TCB,
- 395 2. the TCB executes the *TPM_Quote*, *i.e.*, it signs the PCRs from 0 to 10 and the nonce
 396 using the Attestation Key (AK), and prepares the IMA log for delivery,

- 397 3. the Platform sends back to the Attestor the Quote (*i.e.*, the signature over the values
 398 of the PCR from 0 to 10 and the nonce) and the AK certificate, the PCR values, the
 399 IMA log, and the list of files measured by the bootloader with their digest; then,
 400 4. the Attestor performs all checks to verify the TCB integrity:
- 401 (a) it verifies the validity of the Quote using the public key embedded in the
 402 AK certificate (previously set as trusted), also including the nonce;
 - 403 (b) it verifies that the values of “soft PCRs” calculated from the list of files (and
 404 their digest) measured by the bootloader correspond to the values of PCR 0
 405 to 8 extended as well by the bootloader;
 - 406 (c) it calculates the digest of the AK certificate and the value of a “soft PCR”
 407 and verifies that corresponds to PCR9 value;
 - 408 (d) it verifies the integrity of the IMA log by recalculating the value of a “soft
 409 PCR” from the IMA log itself and comparing it with the value of PCR10;
 - 410 (e) it verifies the files included in the IMA log with the exception of the files
 411 to be excluded, (*e.g.*, the log files) against whitelists/blacklists of digests
 412 previously built and exchanged in a trusted manner. The result of the
 413 integrity status can be: trusted if all files in the IMA log are found in
 414 whitelists; unknown if at least one file in the IMA log is not found in whitelists
 415 but no file is found in blacklists; and untrusted if at least one file in the
 416 IMA log is found in blacklists or if any of the checks from (a) to (d) fails.
- 417 5. Finally, the Attestor closes the TLS channel and takes the appropriate action(s).
 418 Which are the appropriate action(s) depends on the specific Industry 4.0 application and
 419 they are subject to further research.

420 5. Reference Implementation

421 The following paragraphs provide a description of our proposed testbed, intended
 422 to be representative of a typical synchronization distribution network. We have imple-
 423 mented the nodes of the testbed by means of the following components:

- 424 • Raspberry Pi[®] 4 (RPi4) Model B [53], a flexible and high-performance Single Board
 425 Computer with a well-supported set of SW libraries and tools for the Linux envi-
 426 ronment;
- 427 • mosaicHAT [54], an open source hat compatible with RPi4. It is based on the
 428 Septentrio’s mosaic-X5[®] receiver [55], a multi-band, multi-constellation GNSS
 429 module representative of the state of the art and supporting the Galileo signals (*i.e.*,
 430 OSNMA-ready);
- 431 • Infineon OPTIGA[™] TPM SLI 9670 Iridium board [48], an evaluation board with
 432 the widely used TPM2.0 chip.



(a)



(b)

Figure 5. Picture of the Master clock node (a) and Slave clock node (b).

433 In detail, Figure 5 (a) shows a picture of a Master node in our testbed, consisting in
 434 a RPi4 with both the mosaicHAT and the TPM stacked on top of it, while Figure 5 (b)
 435 presents a Slave node (*i.e.*, a RPi4 with the TPM only).

436 These components are instrumental to build a highly flexible and reconfigurable
437 testbed, thus suitable to analyse relevant attack vectors and countermeasures in a controlled and legal framework. Our testbed is also scalable, thus we can easily extend
438 it to emulate different network topologies, potentially including a large number of
439 nodes, but avoiding the complexities related to the operation of real PTP hardware in a
440 synchronization distribution network.
441

442 For this purpose, Figure 6 shows a network topology based on five nodes that we
will adopt in all the tests reported in next section.

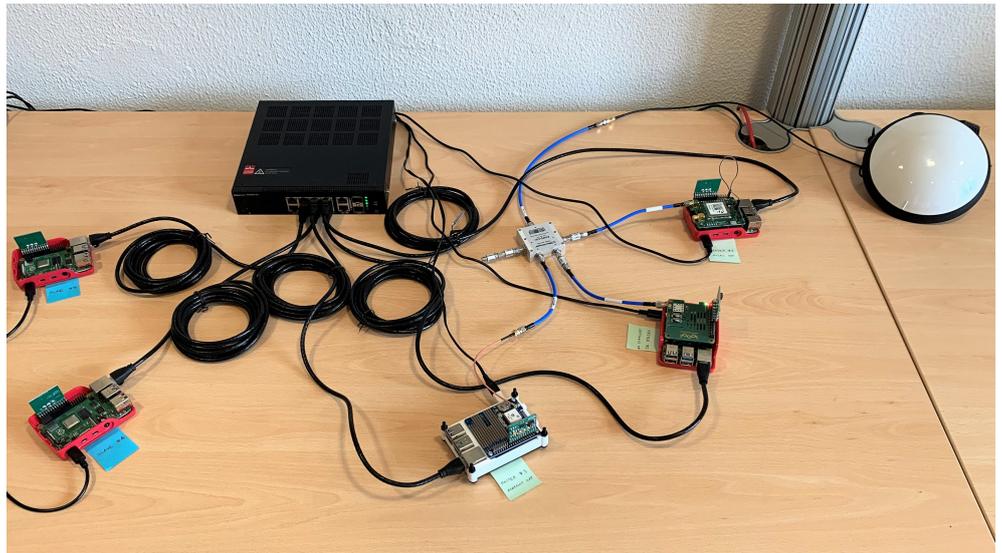


Figure 6. Picture of the full testbed, configured with three Master nodes and two Slave nodes.

443 The picture presents the following components, clockwise starting from the upper
444 right corner:
445

- 446 • a professional GNSS antenna (*i.e.*, Tallysman[®] VSP6037L VeroStar[™] Full GNSS
447 Precision Antenna plus L-band [56]), placed on the desk for illustration purposes
448 only, but properly mounted on the rooftop of the building during the tests;
- 449 • three RPi4 configured as Master nodes (*i.e.*, Master 1, 2, and 3), including both a
450 GNSS module and a TPM stacked on top of it;
- 451 • two RPi4 configured as Slave nodes (*i.e.*, Slave 4 and 5), without a GNSS module.

452 All the five RPi4 nodes in Figure 6 are connected to the power supply and to the
453 same Local Area Network (LAN) by means of a switch and Ethernet cables. As far as
454 the three Master nodes are concerned, we connected them to the same GNSS antenna
455 through a signal splitter. Only the Master 1 is equipped with the mosaicHAT [54], while
456 the other two Master nodes have low cost GNSS modules (*i.e.*, Uputronics[™] Raspberry
457 Pi GPS/RTC Expansion Board [57] for Master 2, Adafruit Ultimate GPS HAT [58] for
458 Master 3). Moreover, the Master 1 is configured as the Grand Master clock for the PTP
459 protocol, while the Master 2 acts as backup master clock and the Master 3 (*i.e.*, the white
460 node in the middle of Figure 6) is used as a monitoring node. In detail, the Master 3 is
461 capable to estimate the relative synchronization errors of the other nodes on the LAN in
462 real time, taking advantage of the ntpq utility [59], and to collect detailed log files, by
463 means of a properly configured ntpd daemon [42].

464 It must be remarked that this setup is potentially suitable to a comparative perfor-
465 mance assessment of the different GNSS modules and/or for emulating different attacks
466 against the GNSS RF Spectrum (see Section 2). Nonetheless, as previously highlighted
467 in Section 1, GNSS RF vulnerabilities are out of the scope for this paper, thus the role of
468 the GNSS modules in our testbed is just to provide a reliable time synchronization to all
469 the RPi4 nodes. For these reasons, the experiments presented in following section will
470 focus only on attacks against the SW integrity and configuration of the nodes.

471 6. Results and Discussion

472 The testbed introduced in previous section is suitable to test different SW stacks
 473 and configurations and to comparatively assess them in terms of performance and
 474 robustness. For instance, we can test both the configurations #1 and #2 as in previous
 475 Figure 2 on the Master node. In the first case, the `gpsmon` utility program [60] allows to
 476 check the behaviour of `gpsd` in real-time [41]. On the other hand, the `ntpq` utility [59] is
 477 usable in both configurations to monitor the status of `ntpd` [42] and then to estimate the
 478 synchronization performance of the Master 1 in our testbed.

479 It is worth to recall that our testbed adopts the `ntpd` daemon to accurately synchro-
 480 nize the local clock of the Master node to the GNSS time scale, while `ptpd` distributes the
 481 timing information from the Master to multiple Slave nodes. Typical time distribution
 482 networks can require a remarkable amount of time (*e.g.*, ranging from few minutes up to
 483 several days) to achieve an initial synchronization between multiple nodes. In fact, small
 484 frequency and phase corrections allow to gradually and continuously adjust (*i.e.*, *disci-*
 485 *pline*) all the clocks. The required amount of time for these operations mainly depends
 486 on the quality of the local oscillators of the nodes, on the specific configuration of `ntpd`
 487 and `ptpd`, and on the application requirements in terms of synchronization accuracy and
 488 precision.

489 As an example, Table 1 summarizes the obtained results running both the previous
 490 configurations #1 and #2 for different time intervals (*i.e.*, from 5 minutes up to 3 days).
 491 In detail, Table 1 reports offset and jitter values (in milliseconds) estimated by `ntpq` for
 492 both 1PPS and NMEA outputs in configuration #1, while in configuration #2 such values
 493 are available just for 1PPS.

Table 1. Obtained time offset and jitter values (in ms) in different configurations

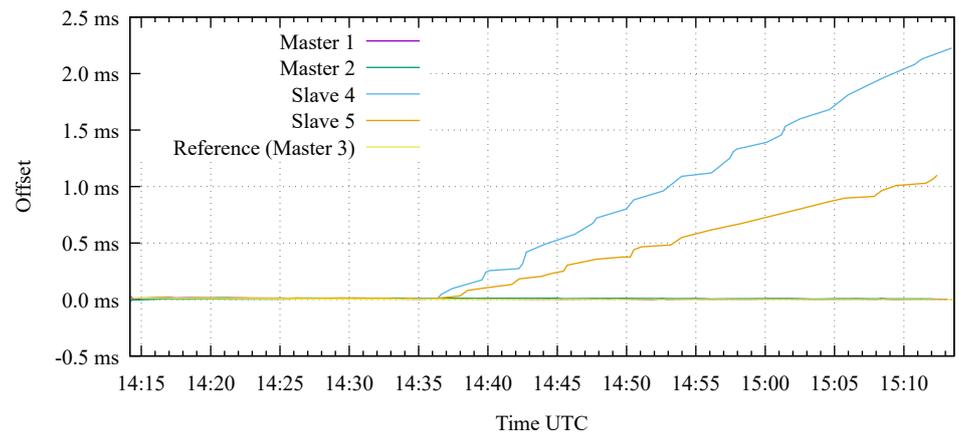
Duration	Obtained outputs Configuration	1PPS		NMEA	
		offset	jitter	offset	jitter
5 minutes	#1	-4.224	1.106	-4.477	1.147
	#2	-3.636	0.896	–	–
10 minutes	#1	0.038	0.022	-0.306	0.135
	#2	0.032	0.018	–	–
1 hour	#1	0.001	0.001	0.190	0.130
	#2	-0.001	0.001	–	–
3 days	#1	0.001	0.001	1.287	0.182
	#2	-0.001	0.001	–	–

494 The 1PPS columns demonstrate that the configurations #1 and #2 result in equivalent
 495 performance after 1 hour (*i.e.*, a synchronization accuracy of the order of 1 μ s). Thus,
 496 we select the configuration #2 in order to simplify the SW stack and to reduce possible
 497 vulnerabilities in the Master nodes and we will use it in the following tests.

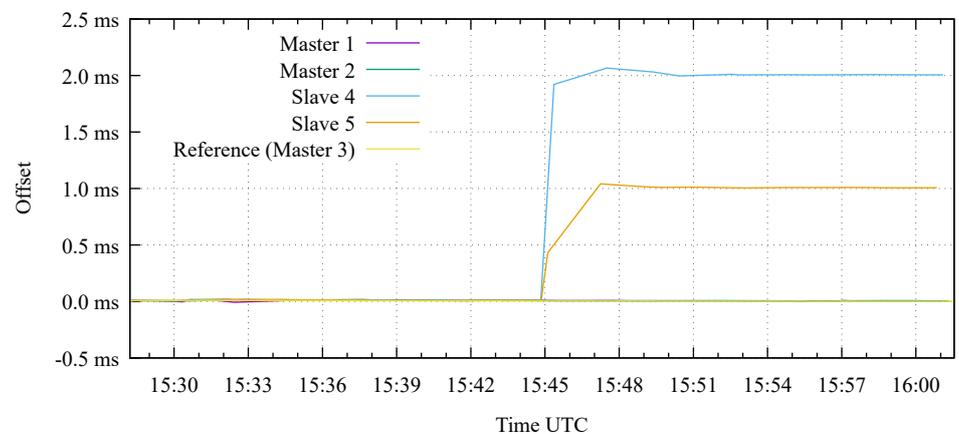
498 The testbed also provides a playground to test different attacks. For instance, the
 499 network topology shown in Figure 6 is suitable to emulate specific attacks against the
 500 PTP SW integrity of the nodes, according to the previous discussion in Section 3.2. In
 501 detail, Figure 7 reports the obtained results considering the following attacks:

- 502 1. an attack against the SW integrity of the `ptpd` daemon;
- 503 2. an attack against the `ptpd` configuration file.

504 The Attack 1 emulates an attacker that has gained access with superuser privileges
 505 to the two Slave nodes. The attack consists in the installation of a maliciously modified
 506 executable file on both the nodes, obtained by modifying the original source code of
 507 `ptpd` [46]. In our test, the modified SW introduces an instantaneous time bias on all the
 508 time stamps received by each Slave node. We insert such bias in the lowest possible level
 509 of the `ptpd` source code, as close as possible to the NIC hardware (*i.e.*, by adding 12 lines
 510 of code to the `net.c` source file). The instantaneous bias has an initial small magnitude,



(a)



(b)

Figure 7. Obtained results on the testbed during the first attack against the SW integrity (a) and the second attack against the configuration of the PTP daemon (b).

511 intended to be undetectable with conventional measures (*e.g.*, on ptpd statistics and
 512 event logs), and it gradually increases with a consistent sign over the duration of the
 513 attack, thus resulting in a frequency drift. In this way, a sufficiently long attack duration
 514 can produce arbitrarily large time offsets on the target nodes, potentially with different
 515 relative time offsets (*i.e.*, different error magnitude and/or sign). As shown in Figure
 516 7(a), the first test lasts for approximately one hour and begins with an initial phase in
 517 nominal conditions, where all the five nodes are correctly synchronized. Then, the attack
 518 actually starts to introduce a time bias at 14:36. In this test, we intentionally introduce a
 519 different frequency drift in the two victim nodes (*i.e.*, $1 \mu\text{s}/\text{s}$ on Slave 4 and $0.5 \mu\text{s}/\text{s}$ on
 520 Slave 5) during the second part of the test. For this reason, the Slave 4 reaches a time
 521 synchronization error larger than 2 ms at the end of the test, while the Slave 5 shows an
 522 error larger than 1 ms with respect to all the other Master nodes.

523 On the other hand, the Attack 2 emulates a simpler scenario where we assume that
 524 the attacker is only capable to edit the ptpd configuration file (and not its executable
 525 file). In practice, the modified configuration file on each target Slave node introduces a
 526 constant time bias with an arbitrarily large magnitude on all the time stamps received
 527 from the Master node. Such large bias can result in a time step on the victim clocks during
 528 the attack execution, potentially detectable with conventional measures (*e.g.*, on ptpd
 529 statistics and event logs). We emulate this attack on our testbed by setting a wrong value
 530 for one of the parameters in the ptpd configuration file (*i.e.*, `ptpengine:offset_shift`).

531 In this way, we introduce an error of 2 ms and 1 ms for Slave 4 and Slave 5, respectively,
532 thus resulting in different relative time offsets, as demonstrated in Figure 7(b).

533 Trusted Computing principles enables the mitigation of both these attacks. In
534 fact, the Remote Attestation is capable to detect the unauthorized replacement of an
535 executable file (or a shared library), when the Attestor compares the IMA Log with
536 whitelists/blacklists. The same consideration applies for an unauthorized modification
537 to a configuration file.

538 In the following example we first measure as good the executable of the ptpd
539 daemon and its configuration file, as the measurements in the IMA log coincide with the
540 ones in whitelist. An excerpt of the output of the command showing the IMA log is as
541 follows:

```
pi@raspberrypi:~/$ sudo cat /sys/kernel/security/ima/ascii_runtime_measurements
10 03eb317e687cbd21e360e257b4810f8ac711fb4c ima 9797edf8d0eed36b1cf92547816051c8af4e45ee boot_aggregate
10 1110984f14fc70c87f90053612c3feaa068f66d6 ima 339c9d9d10dfc25731d6f3d00b80e173e35456c /lib/systemd/systemd
10 bc447ab6e2f476455644dbcf9187c922418f02ba ima 369c4027e9b09131c6971ee963bfd37d41dc251c /lib/arm-linux-gnueabi/ld-2.28.so
10 e359bd4b3a5b64f125dda645958237a123fe9fe7 ima 9e7916b2b9232f7db4cff863a6588e10253c0285 /etc/ld.so.preload
...
10 26bf9166f38fd64c452484ddc48d91c32913b53b ima 77a352d6c2817ef4c6df25512f104e46fcf1d6e0 /usr/local/sbin/ptpd2
10 d52997919ea0fe1cdec53bf2546d3db3076e9f58 ima b526bd78d18255929e2c2d2ccd821c47abc10912 /home/pi/PTPd/ptpd2.slave.conf
...
```

542 The last two rows are related to the ptpd executable file and to its configuration file
543 installed in the Slave 4 and Slave 5 nodes, including their full pathnames.

544 An excerpt of the whitelist file is the following:

```
b526bd78d18255929e2c2d2ccd821c47abc10912 /home/pi/PTPd/ptpd2.slave.conf
77a352d6c2817ef4c6df25512f104e46fcf1d6e0 /usr/local/sbin/ptpd2
01a2ecb8982b43a7e3cd9fb9af2e239def49d073 /usr/local/sbin/tpm2-abrmd
bc595d77c2cea5eb899927d0a26d88292eaa64f5 /usr/local/bin/ima_boot_aggregate
fab6bd62f87f58395c37e470a9c1efd2d6d08f4c /usr/local/bin/ima_measure
a93155873b9f2becbbab6e8641b5a2d566fd678f /usr/local/bin/ima_mmap
20436ac9bbf6eeee3567bcd3fd08b4c845578df0 /usr/local/bin/ra_build_white_list
```

545 Looking at the first two rows of the whitelist, we can appreciate that the digest
546 values for both the configuration file and the ptpd executable are consistent with the
547 corresponding values of the previous IMA log.

548 At this point, we simulate the previous Attack 1 (*i.e.*, an unauthorized SW modifi-
549 cation) by stopping the PTP daemon, replacing the ptpd executable with a maliciously
550 modified version, and then running it with the nominal configuration, as follows:

```
pi@raspberrypi:~/$ sudo killall ptpd2
pi@raspberrypi:~/$ sudo cp ~/PTPd/ptpd/src/ptpd2 /usr/local/sbin/ptpd2
pi@raspberrypi:~/$ sudo ptpd2 -c /home/pi/PTPd/ptpd2.slave.conf
```

551 These commands trigger new IMA measurements, which result in the following records
552 appended to the IMA log:

```
10 2e8a62501ac82d1b51f073d019d5e6d41e5999cf ima e1299122fc1dfdb0707a96bf6b879273278c941a /usr/bin/killall
10 c7f7be7019734b0d664c0e0bc8e407923b50d9b03 ima 19e940258bfba6cc25e66c73bb0d7939d9583a1c /home/pi/PTPd/ptpd/src/ptpd2
10 c57611f6a5a369a0b0f9412dd307f78defc8e725 ima 19e940258bfba6cc25e66c73bb0d7939d9583a1c /usr/local/sbin/ptpd2
```

553 By comparing the IMA log and the whitelist, the attack is immediately detected, as
554 the digest for ptpd in the former is different from the new digest in the latter.

555 At the end of that test, we restore the original ptpd and reboot both the Slave nodes
556 in order to return to the nominal synchronization of all the nodes. Then, we can emulate
557 the Attack 2 (*i.e.*, an unauthorized modification on the configuration file). First we
558 retrieve the Process ID (PID) of the running ptpd daemon by watching the content of
559 its status log file, then we overwrite the original configuration file with a maliciously

560 modified version and, finally, we send the appropriate signal (*i.e.*, SIGHUP) to ptpd in
561 order to force it to reload its configuration file:

```
pi@raspberrypi:~/$ watch -n 1 cat /var/log/ptpd2.status.log
pi@raspberrypi:~/$ sudo cp ~/PTPd/ptpd2.slave_offset.conf ~/PTPd/ptpd2.slave.conf
pi@raspberrypi:~/$ sudo kill -s SIGHUP <PID>
```

562 As in previous attack, these commands trigger new IMA measurements, as follows:

```
10 e78bbb133849e6ec45dd9e2dd2b87b13eaa8ba14 ima bffac84554ed0fc938387d163dae108d26f80341 /usr/bin/watch
10 06f51b1aee4b32d001ab3a4c905530de203812ea ima 3fd479d4a21b4c7a07df154156b69c8be7c2b5a1 /home/pi/PTPd/ptpd2.slave_offset.conf
10 fc49caa91dadee8af0540791f96b9f4f3483b56a ima 1ea5039520c120ab8c77e870c096e7b7d6908ee4 /bin/kill
10 d67b1554ac31101b987b1dd9537e8c45a659a705 ima 3fd479d4a21b4c7a07df154156b69c8be7c2b5a1 /home/pi/PTPd/ptpd2.slave.conf
```

563 In this case the attack is detectable by comparing the digest of the configuration file
564 in the IMA log with the previous value stored in the whitelist.

565 It is worth to clarify that, for both attacks, the detection by the Attestor can have a
566 potential delay with respect to the actual attack initiation. Such delay depends on the
567 given periodicity at which the Attestor contacts the nodes to attest/verify the integrity
568 of their software and configuration: a frequent execution of such Remote Attestation
569 protocol allows a rapid detection of an ongoing attack, but it comes at the price of an
570 increased computational load and network overhead. In this sense, the periodicity of the
571 attestation procedure represents an important design parameter to be tailored to each
572 use case in Industry 4.0 applications, trading-off security versus complexity.

573 7. Conclusions

574 This paper has provided a complete analysis of the possible attack vectors to a GNSS-
575 based Time Distribution Network, shedding the light on the importance of deploying
576 a software integrity architecture in Industry 4.0 applications. We have proposed and
577 demonstrated a Trusted Computing implementation for embedded devices, together
578 with a real-world option to make a flexible and scalable testbed for further exploration
579 of threats and vulnerabilities due to the combination of diverse attack vectors.

580 **Author Contributions:** All authors have contributed equally to this work. All authors have read
581 and agreed to the published version of the manuscript.

582 **Funding:** This work was developed within the ROOT project (www.gnss-root.eu), funded by the
583 European GNSS Agency under the European Union's Horizon 2020 – G.A. n. 101004261.

584 **Institutional Review Board Statement:** Not applicable.

585 **Informed Consent Statement:** Not applicable.

586 **Data Availability Statement:** No data available online. For further query please contact the
587 corresponding author.

588 **Acknowledgments:** The authors would like to thank Gianluca Ramunno for the technical support
589 provided during the initial part of this work, all the partners of the ROOT project for their
590 useful suggestions and, finally, the reviewers whose comments helped to improve and clarify this
591 manuscript.

592 **Conflicts of Interest:** The authors declare no conflict of interest.

593 Appendix A

594 Appendix A.1 Default IMA policy

595 The default IMA policy for the measurement, called `ima_tcb`, is the following:

```
596 dont_measure fsmagic=PROC_SUPER_MAGIC
597 dont_measure fsmagic=SYSFS_MAGIC
598 dont_measure fsmagic=DEBUGFS_MAGIC
```

```
599 dont_measure fsmagic=TMPFS_MAGIC
600 dont_measure fsmagic=SECURITYFS_MAGIC
601 dont_measure fsmagic=SELINUX_MAGIC
602 measure func=BPRM_CHECK
603 measure func=FILE_MMAP mask=MAY_EXEC
604 measure func=PATH_CHECK mask=MAY_READ uid=0
```

```
605 and instructs IMA to measure all files loaded for execution and all files read under the
606 user account root, with the exclusion of all special file systems. Files that get modified
607 during OS run-time are re-measured by IMA and can appear multiple time in the IMA
608 log with different measurements. In our prototype the variable files, like the log files
609 located in /var/log are measured but excluded during the verification that takes place
610 with the Remote Attestation.
```

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Andrea Vesco, received the M.Sc. degree in Telecommunication Engineering and the Ph.D. in Computer and System Engineering from the Politecnico di Torino in 2003 and 2009 respectively. After one year of post-doc with the Control and Computer Engineering Department as a member of the Computer Networks Group at the Politecnico di Torino, he joined the Networking Lab at the Istituto Superiore Mario Boella (ISMB) in 2010. His research interests were on Quality-of-Service (QoS) over packet switched networks and wireless access networks. Moreover he carried out research activities on QoS over Network-on-Chip (NoC) in collaboration with the System-Level Design Group at the Columbia University of the city of New York. He had the opportunity to focus his researches on Digital Smart Cities from 2013 to 2018. Today, he is with the LINKS Foundation a non-profit institution promoting, conducting and strengthening innovation, research processes and the technology transfer. He leads the cybersecurity research team focusing on the security of connected systems such as IoT and critical infrastructures.